

On the automorphism group of the universal homogeneous meet-tree

Itay Kaplan*

Tomasz Rzepecki[†]

Daoud Siniora

Abstract

We show that the countable universal homogeneous meet-tree has a generic automorphism, but it does not have a generic pair of automorphisms.

1 Introduction

A countable structure M has a *generic automorphism* if its automorphism group has a comeagre conjugacy class. This is an important property which has certain implications on the automorphism group G (for example: every element is a commutator, G cannot be written as a proper free product with amalgamation and more; see [Mac11, Proposition 4.2.12]).

A much stronger property is having *ample generics*, which means that for all $n \in \mathbf{N}$, G has a generic tuple of length n or, in other words, G^n contains a comeagre orbit under the action of G by diagonal conjugation. Having ample generics implies in particular the small index property [Mac11, Theorem 5.2.5] — which is desirable because, for instance, it implies that if M is \aleph_0 -categorical, then the automorphism group G (as a pure group) determines M up to bi-interpretability (among all countable \aleph_0 -categorical structures; see e.g. [Mac11, p. 5.2.2]).

This article started as an attempt to find a counterexample to a question of Dugald Macpherson (also appearing in the third author's thesis [Sin17, Chapter 7, Question 10]) which asks whether there is an ultrahomogeneous \aleph_0 -categorical structure admitting ample generics with the strict order property (i.e., defining a partial order with infinite chains). The obvious candidate, $(\mathbf{Q}, <)$, fails: by the works of Hodkinson (unpublished), Truss [Tru07], and the third author (who gave a new proof of this result) [Sin17, Lemma 6.1.1], we know that $\text{Aut}(\mathbf{Q}, <)$ has no generic pair of automorphisms (although $\text{Aut}(\mathbf{Q}, <)$ does have a generic automorphism by [Tru92; KT01]).

(Order-theoretic) trees are partial orders in which for every element a , the set of elements below it is linearly ordered, while meet-trees are trees with a meet function (see Definition 2.10). They are often used as a basis for interesting examples in the realm of NIP unstable theories (in the sense of Shelah's classification theory). For example, in [KS14b; KS14a], trees were the basis for a counterexample to an old conjecture of Shelah regarding the existence of

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indiscernibles in NIP. See also [Sim15, Section 2.3.1]. Finite meet-trees form a Fraïssé class and thus there is a countable universal ultrahomogeneous and \aleph_0 -categorical meet-tree \mathbf{T} which we call the *universal countable dense meet-tree*. We therefore thought that it would be interesting to understand the automorphism group of \mathbf{T} with respect to generic automorphisms.

Our main result is:

Main Theorem. *Let \mathbf{T} be the universal countable dense meet-tree. Then its automorphism group $G = \text{Aut}(\mathbf{T})$ has a generic automorphism, but not a generic pair of automorphisms.*

(See Corollary 3.8 and Theorem 6.11.) ◇

Note that the **Main Theorem** implies that in particular, $\text{Aut}(\mathbf{T})$ does not have ample generics — which, as was mentioned, would imply having the small index property. However, it follows by [DHM89, Theorem 4.1] that $\text{Aut}(\mathbf{T})$ (as well as the analogues for meet-trees of bounded arity, for which we also show the nonexistence of a generic pair) does have the small index property.

We use a criterion for having a generic automorphism that was established by Truss [Tru92, Theorem 2.1] and then improved to a characterisation independently by Ivanov [Iva99, Theorem 1.2] and Kechris and Rosendal [KR07, Theorem 6.2]. Namely, to show that a ultrahomogeneous structure M with age \mathcal{K} has a generic n -tuple of automorphisms, one needs to prove that the class \mathcal{K}_p^n — of pairs (A, \bar{f}) such that $A \in \mathcal{K}$ and \bar{f} is an n -tuple of partial automorphisms of A — has the joint embedding property (JEP) and a version of the amalgamation property which we call the existential amalgamation property (EAP)¹ (see Fact 2.8).

To show EAP in the case of finite meet-trees, we find a cofinal subclass of \mathcal{K}_p^1 consisting of amalgamation bases (which is the aforementioned criterion from [Tru92]). Starting with some member (A, f) , we extend f to a partial automorphism p on a bigger domain B in such a way that (B, p) is an amalgamation base in \mathcal{K}_p^1 . The idea in finding p is model-theoretic: instead of giving a precise description of p , we define it as being “pseudo existentially closed” in the sense that, roughly, any behaviour that happens in some extension of p is already witnessed in p itself (see Definition 6.3). This method eliminates the need for a careful analysis of the interactions between multiple orbits of p that we might otherwise need.

Finally let us remark that while we did not check all the details, our methods seem to recover a proof of existence of generics for $(\mathbf{Q}, <)$ (see Remark 6.13), and also provide a similar result for lexicographically ordered meet-trees (see Remark 6.14). They could also be helpful in finding generics in the case of meet-trees of bounded arity (see Remark 6.12; see also Corollary 3.8 for the nonexistence of generic pairs in this case).

The paper is organized as follows. In Section 2 we recall Fraïssé classes and formally define EAP. We then discuss trees and give their basic properties. In Section 3 we prove that the automorphism group of the countable dense tree does not have a generic pair of automorphisms (in fact we give a more general statement, also about trees with bounded arity). In Section 4 we follow [KT01] and discuss *determined* finite partial automorphisms in an abstract context, giving a sufficient condition under which they form a class of amalgamation bases. In Section 5 we discuss the possible orbits of partial automorphisms in meet-trees. Finally, in Section 6 we prove that the class of determined partial automorphisms of meet-trees is indeed cofinal, thus proving EAP for the class of finite meet-trees.

¹In [KR07] and [Iva99] this property was called the “weak amalgamation property” (WAP), and “almost amalgamation property”, respectively, but we chose this term, coined in an unpublished work of Ben-Yaacov and Melleray, as it is more descriptive.

We end the introduction with some open questions.

Our main results are similar to the ones in [KT01; Tru07] regarding $(\mathbf{Q}, <)$ and the universal partial order (although it is not known if the universal partial order has a generic pair); more recently, a new preprint [KM19] appeared giving similar results on two different structures (the universal ordered boron tree — roughly speaking, a graph theoretic binary tree with a lexicographical order — and the universal ordered poset). The latter’s motivation came from a different yet related question of finding an ultrahomogeneous *ordered* structure whose automorphism group has ample generics and is extremely amenable (in other words, by [KPT05], its age has the Ramsey property).

Even more recently, a preprint [Duc19] appeared, which analyses the homeomorphism groups of the so-called Ważewski dendrites, in particular showing the analogue of the **Main Theorem** for the group of homeomorphisms of the Ważewski dendrite $D_{\{\infty\}}$ (into which $\text{Aut}(\mathbf{T})$ naturally embeds as a meagre subgroup, namely as the stabiliser of an end point). We were not aware of this work. By private communication with Duchesne, it appears that though it is not written explicitly there, our result follows from the main theorem there and vice versa. Both papers use the aforementioned criterion for establishing the existence of a generic automorphism, but the methods of finding amalgamation bases are different. (Interestingly, from the point of view of [Duc19], meet-trees with bounded arity are analogues to dendrites with bounded branching. The latter do not have a dense conjugacy class according to [Duc19, Proposition 1.1], while the former do, see also Remark 6.12.)

Question 1.1. What is the correct generalization of all these results?

Model theoretically, it seems appealing to consider the situation in general NIP \aleph_0 -categorical structures with perhaps some further restrictions, as was done in [Sim18]. This situation does not quite generalize ours since trees are not “rank 1” in the sense defined there, and the universal partial order is not NIP.

Another natural question that comes to mind is the following. In all the examples known to us, if EAP occurs for \mathcal{K}_p^1 then in fact there is a cofinal class of amalgamation bases (a property we denote by CAP, see Definition 2.6).

Question 1.2. Is it always the case that EAP is equivalent to CAP for the class \mathcal{K}_p^1 (where \mathcal{K} is any Fraïssé class)?

We are not sure what the situation is with trees of bounded arity; it may be a candidate for a counterexample — see Remark 6.12.

Since we know that the generic automorphism exists, it seems natural to ask how complicated it is.

Question 1.3. What can be said, model theoretically — in terms of classification theory — about the structure (\mathbf{T}, σ) , where σ is a generic automorphism?

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2 Preliminaries

Fraïssé classes and limits

We briefly recall the basic notions related to Fraïssé classes. See [Hod97, Chapter 6] for more exposition. In contrast to [Hod97], it will be convenient for us to consider classes of structures with partial functions. Formally, they can be thought of as relations (via their graphs), so this is only a superficial change.

Definition 2.1. Let \mathcal{K} be a class of first order structures (possibly with some partial functions), closed under isomorphisms.

- We say that \mathcal{K} has the *hereditary property (HP)*, if given any $A \in \mathcal{K}$ and substructure $B \subseteq A$, we have $B \in \mathcal{K}$.
- We say that \mathcal{K} has the *joint embedding property (JEP)* if given any $A_1, A_2 \in \mathcal{K}$, there is some $B \in \mathcal{K}$ such that both A_1 and A_2 can be embedded in B .
- We say that $A \in \mathcal{K}$ is an *amalgamation base* (in \mathcal{K}) if given any $B_1, B_2 \in \mathcal{K}$ and embeddings $i_j: A \rightarrow B_j$ for $j = 1, 2$, there is some $C \in \mathcal{K}$ and embeddings $i'_j: B_j \rightarrow C$ for $j = 1, 2$ such that $i'_1 \circ i_1 = i'_2 \circ i_2$.
- We say that \mathcal{K} has the *amalgamation property (AP)* if every $A \in \mathcal{K}$ is an amalgamation base in \mathcal{K} .
- We say that \mathcal{K} is *uniformly locally finite* if given any $n \in \mathbf{N}$, there is an upper bound on the size of an n -generated element of \mathcal{K} .
- A *Fraïssé class* is a class of finitely generated structures which is closed under isomorphisms, and has HP, JEP and AP.
- The *age* $\text{Age}(M)$ of a first-order structure M is the class of all (isomorphism types of) finitely-generated substructures of M . \diamond

We recall the notion of a partial automorphism (which is fundamental for this paper).

Definition 2.2. Given a first order structure M , a partial function $M \rightarrow M$ is called a *partial automorphism* if it preserves the quantifier-free types over the empty set.

Fact 2.3. *If \mathcal{K} is a Fraïssé class, then there is a unique (up to isomorphism), countable structure \mathbf{K} whose age is exactly \mathcal{K} and which is ultrahomogeneous (i.e. every finite partial automorphism of \mathbf{K} extends to an automorphism).*

Furthermore, if \mathcal{K} is uniformly locally finite, then the theory of \mathbf{K} has quantifier elimination and (if \mathbf{K} is infinite) is \aleph_0 -categorical.

Proof. This is classical; see for instance [Hod97, Theorem 6.1.2] and [Hod97, Theorem 6.4.1]. \square

Definition 2.4. The structure \mathbf{K} as in Fact 2.3 is called the (*Fraïssé*) *limit* of \mathcal{K} . \diamond

Partial automorphisms, generic automorphisms

If M is a countable first-order structure, then $\text{Aut}(M)$ has a natural Polish group structure (with the pointwise convergence topology), and we can use descriptive set theory to study it. The notion of a generic is due to [Tru92] and [KR07] (for tuples).

Definition 2.5.

- An element $\sigma \in \text{Aut}(M)$ is called *generic* if its conjugacy class is comeagre in $\text{Aut}(M)$ (i.e. it contains a dense G_δ set).
- More generally, a tuple $(\sigma_1, \dots, \sigma_n) \in \text{Aut}(M)$ is *generic* if its diagonal conjugacy class (i.e. the orbit under the action $\text{Aut}(M)$ on $\text{Aut}(M)^n$ by coordinatewise conjugation) is comeagre.
- We say that $\text{Aut}(M)$ has *ample generics* if it has generic tuples of elements of arbitrary length. \diamond

The definition of EAP below is due to [Iva99] (where it is called *almost amalgamation property*). It is also used in [KR07] (where it is called *weak amalgamation property*).

Definition 2.6. Fix a class \mathcal{K} of first order structures, closed under isomorphisms.

- We say that \mathcal{K} has *EAP* (*existential amalgamation property*) if for every $A \in \mathcal{K}$, there is some $B \in \mathcal{K}$ and an embedding $i_{AB}: A \rightarrow B$ such that for any embeddings $i_{BC}: B \rightarrow C$, $i_{BD}: B \rightarrow D$ (with $C, D \in \mathcal{K}$), there are embeddings $i_{CE}: C \rightarrow E$ and $i_{DE}: D \rightarrow E$ (where $E \in \mathcal{K}$) such that $i_{CE} \circ i_{BC} \circ i_{AB} = i_{DE} \circ i_{BD} \circ i_{AB}$.
- We say that \mathcal{K} has *CAP* (*cofinal amalgamation property*) if we can choose the B such that $i_{CE} \circ i_{BC} = i_{DE} \circ i_{BD}$ (i.e. if \mathcal{K} has a cofinal subclass of amalgamation bases). \diamond

Remark 2.7. If \mathcal{K} has CAP, it has EAP. \diamond

Fact 2.8. Fix any $n \in \mathbf{N}$. If \mathcal{K} is a Fraïssé class with limit \mathbf{K} and \mathcal{K}_p^n is the class of \mathcal{K} -structures with n -tuples of partial automorphisms, then the following are equivalent:

- $\text{Aut}(\mathbf{K})$ has a generic n -tuple,
- \mathcal{K}_p^n has JEP and EAP.

Proof. This is [KR07, Theorem 6.2]. Under the additional assumption that \mathbf{K} is \aleph_0 -categorical (which includes our applications), this is a special case of [Iva99, Theorem 1.2]. \square

Corollary 2.9. If \mathcal{K} is a Fraïssé class with limit \mathbf{K} and \mathcal{K}_p^n has JEP and CAP, then \mathbf{K} has a generic n -tuple of automorphisms.

Proof. Immediate by Remark 2.7 and Fact 2.8. (For $n = 1$ this is essentially [Tru92, Theorem 2.1].) \square

Trees

Definition 2.10.

- A *tree* is a partially ordered set (A, \leq) which is *semilinear* (that is, for every $a_0 \in A$, the set $A_{\leq a_0} = \{a \in A \mid a \leq a_0\}$ is linearly ordered) and such that every pair of elements has a common lower bound.
- A *meet-tree* (or \wedge -tree) (A, \leq, \wedge) is a tree which is also a lower semilattice, i.e. a tree (A, \leq) together with a binary (meet or infimum) function $\wedge : A^2 \rightarrow A$ such that for every $a, b \in A$, $a \wedge b$ is the largest element of $A_{\leq a} \cap A_{\leq b}$. \diamond

Remark 2.11. If (A, \leq) is a tree with the property that every pair has an infimum, then there is a unique way to expand it to a meet-tree. In particular, every finite tree has a unique meet-tree structure. However, not every embedding of finite trees yields an embedding of the resulting meet-trees. \diamond

Remark 2.12. The \wedge operation is associative, commutative and idempotent. \diamond

Definition 2.13. Given a tree T , the *arity* of T is the maximal size of a set $A \subseteq T$ of pairwise incomparable elements such that if $a_1, a_2, a_3 \in A$ are distinct and $b \in T$ is such that $b < a_1$ and $b < a_2$, then $b < a_3$ (or ∞ if there is no finite bound).

Remark 2.14. If T is a meet-tree, arity can be equivalently defined in the following way: we say that T is k -ary if k is the maximal size of a subset $A \subseteq T$ such that all pairs have the same meet, which is not equal to any of them. Note that this definition shows that being at most k -ary is an universal property in the language of meet-trees (it is not hard to see that it is not an universal property in pure order language). \diamond

Fact 2.15. *The class of all finite meet-trees is a Fraïssé class (in the language of meet-trees). Given any positive integer k , the class of all finite k -ary meet-trees is a Fraïssé class (in the language of meet-trees).*

Consequently, there is a countable generic meet-tree, \mathbf{T} , and for every k there is a countable generic k -ary meet-tree, \mathbf{T}_k . \mathbf{T} and each \mathbf{T}_k is \aleph_0 -categorical, ultrahomogeneous and has elimination of quantifiers.

Proof. The first part is straightforward. The second part follows from Fact 2.3. \square

(Notice that in particular, a 1-ary meet-tree is simply linear, and \mathbf{T}_1 is (interdefinable with) the universal linear ordering, isomorphic to (\mathbf{Q}, \leq) .)

We will use the notation \mathbf{T} and \mathbf{T}_k throughout the paper.

Remark 2.16. The class of all finite meet-trees in pure order language is *not* a Fraïssé class — it does not have the amalgamation property; it does, however have a model companion: there is a unique countable existentially closed meet-tree (in pure order language). It is \aleph_0 -categorical and binary, but in contrast to \mathbf{T}_2 , which is a meet-tree, no two incomparable elements have a meet. (See [BBPP18] for more details.) It might be interesting to ask whether it has a generic automorphism, but the methods based on Fact 2.8 used in this paper do not seem to apply directly. \diamond

Remark 2.17. Another approach to trees is by graph theory: we can identify every finite tree with an acyclic directed graph, but the Fraïssé limits of classes of finite trees in this language will be quite different (for instance, the “order” on each branch will not be dense). In [KM19], the authors study the existence of generic automorphisms in this context. \diamond

The following simple observation, reminiscent of the ultrametric triangle inequality, will be immensely useful in the rest of this paper.

Fact 2.18. *Let a, b, c be elements of a meet-tree T . Then:*

- $a \wedge b \wedge c = a \wedge b$ or $a \wedge b \wedge c = a \wedge c$,
- if $a \wedge b > a \wedge c$, then $a \wedge c = b \wedge c$,
- if $a \wedge b \geq a \wedge c$, then $a \wedge c \leq b \wedge c$.

Proof. The proof is straightforward (using semilinearity). \square

Definition 2.19. Given a partial order (P, \leq) , and an $A \subseteq P$,

- a (*lower*) *cut* in A is a downwards closed, (upwards) directed subset of A (i.e. a $C \subseteq A$ such that for any $c \in C$ and $a \in A$, if $a \leq c$, then $a \in C$, and for any $c_1, c_2 \in C$, there is some $c \in C$ such that $c \geq c_1, c_2$).
- an *order type* over an A is simply a complete quantifier-free type over A in P , in pure order language (with equality). \diamond

Remark 2.20. Fix a partial order (P, \leq) (possibly with some additional structure) and a subset $A \subseteq P$.

Given an order type p over A , we have two corresponding cuts in A , namely $p_{\geq} := \{a \in A \mid p \vdash x \geq a\}$ and $p_{>} := \{a \in A \mid p \vdash x > a\}$. The two are equal if and only if p is not realised in A . If P is linear, then the two cuts uniquely determine p , but in general, it is not true.

Conversely, given a *nonempty* cut $C \subseteq A$, there is an order type p over A such that $p_{\geq} = C$. \diamond

Definition 2.21. Given any $b \in P$, by the *order type of b over A* , $\text{otp}(b/A)$, we mean simply the quantifier-free type of b over A in the order language, and by *the cut of b in A* we mean simply the cut $\{a \in A \mid b \geq a\}$ (which is the same as $\text{otp}(b/A)_{\geq}$ in the notation of Remark 2.20). \diamond

Remark 2.22. In a tree, a directed set is linear, so a cut is simply a downwards closed chain. \diamond

Remark 2.23. Note that for any poset (P, \leq) , the cuts in P are ordered simply by inclusion. If we denote by \hat{P} the set of all cuts in P , partially ordered by inclusion, it is easy to see that:

- P naturally embeds into \hat{P} (where each element is identified with its cut in P);
- \hat{P} is complete in the sense that every directed subset of \hat{P} has a supremum (namely, the union);
- if P is a tree, then so is \hat{P} , and moreover, \hat{P} has a canonical meet-tree structure, with the meet given by intersection;
- if (P, \leq, \wedge) is a meet-tree, then (P, \leq, \wedge) is a substructure of $(\hat{P}, \subseteq, \cap)$. \diamond

The following fact appears to be folklore.

Fact 2.24. Let B be a meet-tree, $b \in B$, and $A \subseteq B$ a substructure (in the language of meet-trees). Put $b' = \max\{x \wedge b : x \in A\}$, and let $a \in A$ be such that $a \geq b'$. Then the quantifier-free type $\text{qftp}(b/A)$ is determined by knowing whether $b = b'$ or $b' < b$, the value of a , and the order type of b' over $A_{\leq a}$.

(In particular, if B has quantifier elimination (for instance, if $B = \mathbf{T}$ or $B = \mathbf{T}_k$ for some $k > 0$, in the sense of Fact 2.15), this also determines $\text{tp}(b/A)$.)

Proof. The proof is left as an exercise to the reader. \square

3 There is no generic pair

In this section, we will use Fact 2.8 to show that the countable generic meet-trees \mathbf{T} and \mathbf{T}_k do not admit generic pairs of automorphisms.

Definition 3.1.

- Given a pair (g_1, g_2) of partial functions, we say that (g'_1, g'_2) is an *extension* of (g_1, g_2) if g'_1, g'_2 are partial functions such that $g_1 \subseteq g'_1$ and $g_2 \subseteq g'_2$.

- Given two pairs $(g'_1, g'_2), (g''_1, g''_2)$ of partial automorphisms of a partially ordered set P and a point $a \in P$, we say that they are *irreconcilable over a* if there is no partial ordering Q and partial automorphisms f_1, f_2 of Q such that $(P, g'_1, g'_2, a), (P, g''_1, g''_2, a) \hookrightarrow (Q, f_1, f_2, b)$ for some $b \in Q$. \diamond

Remark 3.2. If $(g'_1, g'_2), (g''_1, g''_2), (g'''_1, g'''_2)$ are pairs of partial automorphisms of a poset P such that (g'_1, g'_2) and (g''_1, g''_2) are irreconcilable over $a \in P$, while (g'''_1, g'''_2) extends (g'_1, g'_2) , then it is easy to see that:

- (g'''_1, g'''_2) and (g''_1, g''_2) are irreconcilable over a ,
- there is no $f \in \text{Aut}(P)$ fixing a , such that $g'_1 \cup fg''_1 f^{-1}$ and $g'_2 \cup fg''_2 f^{-1}$ are partial automorphisms. \diamond

Proposition 3.3. *Suppose L is a dense linear order, unbounded from below. Suppose g is a finite partial automorphism of L , $a, b \in L$ are such that $b \leq g(a) \leq a$. Then there are $d_1, d_2 \in L$ such that $g \cup \{(d_1, b), (b, d_2)\}$ is a partial automorphism.*

Proof. We may assume without loss of generality that L is countable (by replacing it with a countable dense subset containing b , the range and the domain of g). In this case, there is some $K \supseteq L$, countable, totally ordered without endpoints; since L is dense without lower bound, we may assume that for every $k \in K \setminus L$, we have $k > a$.

Countable dense orderings without endpoints are ultrahomogeneous, so g can be extended to some $f \in \text{Aut}(K)$. Put $d_1 := f^{-1}(b), d_2 := f(b)$. Now, we have $b \leq a$, so $d_2 = f(b) \leq f(a) = g(a) \leq a$, so $d_2 \in L$. On the other hand, $b \leq g(a) = f(a)$, so $d_1 = f^{-1}(b) \leq f^{-1}(f(a)) = a$, so $d_1 \in L$. Since f is a partial automorphism, so is $g \cup \{(d_1, b), (b, d_2)\} \subseteq f$, so we are done. \square

The following fact is essentially [Sin17, Lemma 6.1.1], but we slightly strengthen the conclusion using Proposition 3.3, and we give a more detailed proof.

Proposition 3.4. *Let $(L, <)$ be a dense linear order, unbounded from below, and take some $a \in L$. Let q_1, q_2 be finite partial automorphisms of L such that $q_1(a), q_2(a) < a$. Then (q_1, q_2) admits two irreconcilable extensions (q'_1, q'_2) and (q''_1, q''_2) to pairs of partial automorphisms of L .*

Proof. Given a finite partial automorphism g of L and an element $b_0 \in L$, put $C(b_0, g) := \#\{b \in \text{dom}(g) \mid b \leq b_0\} + \#\{b \in \text{range}(g) \mid b \leq b_0\}$. Note that $C(b_0, g)$ is always a non-negative integer; note also that C is monotone in b_0 : if $b'_0 \leq b_0$, then $C(b'_0, g) \leq C(b_0, g)$. Given a pair (g_1, g_2) of finite partial automorphisms, denote by $m(g_1, g_2)$ the minimal element of the $\langle g_1, g_2 \rangle$ -orbit of a (i.e. the smallest element that can be obtained from a by successive applications of $g_1, g_2, g_1^{-1}, g_2^{-1}$). Write $C_m(g_1, g_2)$ for $C(m(g_1, g_2), g_1) + C(m(g_1, g_2), g_2)$.

Call a pair (g_1, g_2) *minimal* if $C_m(g_1, g_2)$ is minimal among its extensions. We may assume without loss of generality that (q_1, q_2) is minimal (otherwise, we can simply extend it).

Write c for $m(q_1, q_2)$ and B for the union of domains and ranges of q_1 and q_2 . Clearly, $c \in B$. Furthermore, $c \leq q_1(a), q_2(a)$, so it satisfies the assumptions of Proposition 3.3 for both q_1 and q_2 .

We claim that c is in only one of $\text{dom}(q_1), \text{dom}(q_2), \text{range}(q_1), \text{range}(q_2)$. We will show that $c \notin \text{range}(q_1) \cap \text{range}(q_2)$ (the other cases are either analogous or easy to see). Suppose this is not the case. Since c is minimal in its orbit (and $a > c$ is in its orbit), either $q_1^{-1}(c) > c$ or $q_2^{-1}(c) > c$. Suppose without loss of generality that the former holds. Let $d \in L$ be such that

$g_1 := q_1 \cup \{(c, d)\}$ is a partial automorphism (which we have by Proposition 3.3). Then, since $c < q_1^{-1}(c)$, we have $d < c$. It is easy to see that $C(d, q_2) \leq C(c, q_2) - 1$ (we do not count c on the left-hand side) and

$$C(d, g_1) \leq C(d, q_1) + 1 \leq C(c, q_1) - 1 + 1 = C(c, q_1).$$

(The first inequality is that the only extra element we count in $C(d, g_1)$ is d , the second one is because c is counted in $C(c, q_1)$, but not in $C(d, q_1)$.) Since $m(g_1, q_2) \leq d$, we conclude that $C(m(g_1, q_2), g_1) + C(m(g_1, q_2), q_2) < C(c, q_1) + C(c, q_2)$, contradicting minimality of (q_1, q_2) .

From now, we consider the case when $c \in \text{range}(q_1)$ (the other cases are analogous), whence $c \notin \text{dom}(q_2) \cup \text{range}(q_2)$.

Let $c^+ > c$ be such that $(c, c^+] \cap B = \emptyset$, and let $c^- < c$ be such that $[c^-, c) \cap B = \emptyset$. Note that c^+ exists by density and the fact that $a > c$, while c^- exists by density and the assumption that L has no lower bound. We claim that $q_2' := q_2 \cup \{(c, c^+)\}$ and $q_2'' := q_2 \cup \{(c, c^-)\}$ are both partial automorphisms of L . Then if we put $q_1' = q_1'' = q_1$, then clearly (q_1', q_2') and (q_1'', q_2'') will be irreconcilable over a . Note that since $c \notin \text{dom}(q_2)$, we already know that q_2' and q_2'' are well-defined partial functions, and by choice of c^- and c^+ , they are both injective, so it is enough to show that they preserve the order.

Take any $b \geq c$, $b \in \text{dom}(q_2)$. We need to show that $q_2(b) \geq c^+, c^-$. We know that $q_2(b) \neq c$ (because $c \notin \text{range}(q_2)$) and of course $q_2(b) \in B$, so we have $q_2(b) \notin [c^-, c^+]$. It follows that it is enough to show that $q_2(b) \geq c^-$. Suppose towards contradiction that $q_2(b) < c^-$. Let $d \in L$ be such that $g_2 := q_2 \cup \{(c, d)\}$ is a partial automorphism of L (which we have by Proposition 3.3). Since we have $b \geq c$, we have $d \leq q_2(b) < c^- < c$, so $d < c$. As before, we have $C(d, q_1) \leq C(c, q_1) - 1$ and $C(d, g_2) \leq C(d, q_2) + 1 \leq C(c, q_2) - 1 + 1$. Clearly, $m(q_1, g_2) \leq d$, so we have $C_m(q_1, g_2) \leq C(d, q_1) + C(d, g_2) < C(c, q_1) + C(c, q_2)$, a contradiction.

Now, given some $b \leq c$ with $b \in \text{dom}(q_2)$, we need to show that $q_2(b) \leq c^+, c^-$. As before, it is enough to show that $q_2(b) \leq c^+$. Arguing by contradiction, as in the preceding paragraph, using Proposition 3.3, we can find $d < c$ such that $g_2 := q_2 \cup \{(d, c)\}$ is a partial automorphism, leading to a contradiction of minimality of (q_1, q_2) . \square

Lemma 3.5. *Let T be a meet-tree and take an arbitrary $a \in T$. Suppose f, g are partial automorphisms of T such $\text{dom}(f)$ and $\text{dom}(g)$ are closed under \wedge , and satisfy the condition that for every $\eta \in \text{dom}(f)$, there is some $a_\eta \in \text{dom}(g)$ such that $a_\eta \geq \eta$ and $g|_{T_{\leq a_\eta}} \subseteq f$.*

Then $f \cup g$ is a partial automorphism of T .

Proof. Put $h := f \cup g$. Note that if $\eta \in \text{dom}(f) \cap \text{dom}(g)$, then $g|_{T_{\leq a_\eta}} \subseteq f$, so in particular, $f(\eta) = g(\eta)$, so h is a well-defined partial function. Note also that given $\eta \in \text{dom}(f)$, we have $h|_{T_{\leq a_\eta}} \subseteq f$.

Claim. *The domain $\text{dom}(h)$ is closed under \wedge .*

Proof. Since $\text{dom}(f)$ and $\text{dom}(g)$ are closed under \wedge , it is enough to show that if $\eta \in \text{dom}(f)$ and $\nu \in \text{dom}(g)$, then $\eta \wedge \nu \in \text{dom}(h)$.

If $\eta \wedge \nu = \eta$, then there is nothing to prove. Otherwise, $\eta \leq a_\eta \in \text{dom}(g)$, so $a_\eta \wedge \eta = \eta > \eta \wedge \nu$. By Fact 2.18, we infer that $\eta \wedge \nu = a_\eta \wedge \nu$; since $a_\eta \in \text{dom}(g)$, the conclusion follows. \square (claim)

We will show that h is an isomorphism between substructures of T . To that end, we need to show that h is injective and that for any $\eta, \nu \in \text{dom}(h)$, we have $\eta \leq \nu$ if and only if $h(\eta) \leq h(\nu)$, and that $h(\eta \wedge \nu) = h(\eta) \wedge h(\nu)$.

Notice that f^{-1} and g^{-1} satisfy the assumptions of the proposition we are proving: indeed, given $\eta' = f(\eta) \in \text{dom}(f^{-1})$, we have some $a_\eta \in \text{dom}(g)$ with $a_\eta \geq \eta$ and $g \upharpoonright_{T_{\leq a_\eta}} \subseteq f$. But this clearly implies that $f(a_\eta) = g(a_\eta) \geq \eta'$ and $g^{-1} \upharpoonright_{T_{\leq f(a_\eta)}} \subseteq f^{-1}$. Thus by the first paragraph of this proof, $h^{-1} = f^{-1} \cup g^{-1}$ is a well-defined partial function, so h is injective. By the same token, it is enough to show that h preserves meets and that $\eta \leq \nu$ implies $h(\eta) \leq h(\nu)$ (i.e. the converse will follow).

Now, let us fix arbitrary $\eta, \nu \in \text{dom}(h)$ such that $\eta \leq \nu$. We need to show that $h(\eta) \leq h(\nu)$. If $\nu \in \text{dom}(f)$, then it easily follows that $\eta \in \text{dom}(f)$, so it is enough to consider the case when $\eta \in \text{dom}(f)$ and $\nu \in \text{dom}(g)$. It follows that $a_\eta \wedge \nu \in \text{dom}(g)$, so in fact $a_\eta \wedge \nu \in \text{dom}(f) \cap \text{dom}(g)$ (because $a_\eta \wedge \nu \leq a_\eta$); since trivially $a_\eta \wedge \nu \geq \eta$, we have:

$$h(\nu) = g(\nu) \geq g(a_\eta \wedge \nu) = h(a_\eta \wedge \nu) = f(a_\eta \wedge \nu) \geq f(\eta) = h(\eta).$$

Now, we need to show that $h(\eta \wedge \nu) = h(\eta) \wedge h(\nu)$. We may assume that $\eta \in \text{dom}(f)$, $\nu \in \text{dom}(g)$, and also that $\eta > \eta \wedge \nu$ (if we have equality, then $\eta \leq \nu$, and the conclusion follows from the preceding paragraph). Then $\eta = a_\eta \wedge \eta > \eta \wedge \nu$, so (by Fact 2.18) $\eta \wedge \nu = a_\eta \wedge \nu$. As in the preceding paragraph, we have that $a_\eta \wedge \nu \in \text{dom}(f) \cap \text{dom}(g)$. Thus, we have $f(a_\eta) \wedge f(\eta) = f(\eta) > f(\eta \wedge \nu)$. On the other hand $g(\eta \wedge \nu) = g(a_\eta \wedge \nu) = g(a_\eta) \wedge g(\nu)$. In conclusion, we have $h(a_\eta) \wedge h(\eta) > h(\eta \wedge \nu) = h(a_\eta) \wedge h(\nu)$. It follows by Fact 2.18 that $h(\eta \wedge \nu) = h(a_\eta) \wedge h(\nu) = h(\eta) \wedge h(\nu)$, so we are done. \square

Corollary 3.6. *Suppose p_1, p_2 are finite partial automorphisms of a dense, unrooted meet-tree M , such that for some $a \in M$, we have $p_1(a) = p_2(a) < a$. Then (p_1, p_2) admits extensions (p'_1, p'_2) and (p''_1, p''_2) which are irreconcilable over a .*

Proof. Note that since M is dense and unrooted, $M_{\leq a}$ is a dense linear ordering, without lower bound.

For $i = 1, 2$, write $q_i := p_i \upharpoonright_{M_{\leq a}}$, and let q'_i, q''_i be the extensions given by Proposition 3.4 (to partial automorphisms of $M_{\leq a}$), and put $p'_i := p_i \cup q'_i$, $p''_i := p_i \cup q''_i$. Then by Lemma 3.5 (with $f = q'_i$ or $f = q''_i$, $g = p_i$ and $a_\eta := a$ for all η), p'_i and p''_i are partial automorphisms, and clearly, (p'_1, p'_2) and (p''_1, p''_2) are irreconcilable over a . \square

Corollary 3.7. *If M is a dense and unrooted meet-tree and $\mathcal{K} := \text{Age}(M)$, then \mathcal{K}_p^2 (the class of \mathcal{K} -structures with pairs of partial automorphisms) does not have the EAP (see Definition 2.6).*

Proof. Since M is unrooted, in particular, it contains two elements $a > b$ and $p_1^0 = p_2^0 = \{(a, b)\}$ is a partial automorphism. Then any extension (p_1, p_2) of (p_1^0, p_2^0) satisfies the hypothesis of Corollary 3.6, so it admits two extensions irreconcilable over a . This clearly implies the failure of EAP. \square

The following corollary is the second half of the **Main Theorem** (the first half we will prove later in Theorem 6.11).

Corollary 3.8. *If \mathbf{K} is one of \mathbf{T} or \mathbf{T}_k for $k > 0$ (in particular, if it is the dense linear ordering), then \mathbf{K} does not admit a generic pair of automorphisms.*

Proof. In each case, $\mathcal{K} := \text{Age}(\mathbf{K})$ is a Fraïssé class with limit \mathbf{K} . By Fact 2.8 and Corollary 3.7, it follows that the limit of \mathcal{K} (i.e. \mathbf{K}) does not admit a generic pair of automorphisms. \square

4 Determined partial automorphisms

We aim to show that the universal countable meet-tree admits a generic automorphism (even though, by Corollary 3.8, we already know it does not admit a generic pair of automorphisms). Very broadly, the proof follows [KT01]. More precisely, we will find a sufficient condition for a partial automorphism to be an amalgamation base in the class \mathcal{K}_p^1 (where \mathcal{K} is the class of finite meet-trees), and in the next section, we will find a cofinal class of automorphisms satisfying this condition, thus showing CAP for \mathcal{K}_p^1 . This, in conjunction with Corollary 2.9, will give us the existence of generics.

Determined partial automorphisms in an abstract context

The notion of a strict extension and a determined automorphism is due to [KT01]. We have slightly modified it: the authors of [KT01] do not ask that the domain of a strict extension is a substructure, which is a trivial requirement in the case of relational structures. We also introduce the notion of a strictly positive extension.

Definition 4.1. Let M be a first order structure, and let p be a finite partial automorphism of M .

- We say that an extension $f \supseteq p$ of partial automorphisms of M is *strict* if it is an automorphism of a substructure (i.e. $\text{dom}(f) = \text{range}(f)$ is a substructure of M) and $\text{dom}(f)$ is generated by the f -orbits of elements of $\text{dom}(p)$.
- We say that an extension $f \supseteq p$ of partial automorphisms of M is *positively strict* if f is an endomorphism of a substructure (i.e. $\text{dom}(f) \supseteq \text{range}(f)$ and $\text{dom}(f)$ is a substructure of M) and $\text{dom}(f)$ is generated by the *positive* f -orbits of elements of $\text{dom}(p)$ (i.e. the images of $\text{dom}(p)$ by the positive powers of f).
- Given two [positively] strict extensions f_1, f_2 of p , we say that f_1 and f_2 are *isomorphic over p* if there is an isomorphism $\theta: \text{dom}(f_1) \rightarrow \text{dom}(f_2)$ fixing $\text{dom}(p)$ pointwise, such that $\theta \circ f_1 = f_2 \circ \theta$ (note that we do *not* require that θ extends to an automorphism of M).
- We say that p is [*positively*] *determined* if, up to isomorphism over p , it admits a unique [*positively*] strict extension. \diamond

Recall that we want to find a cofinal class of amalgamation bases in \mathcal{K}_p^1 . Proposition 4.2 and Lemma 4.3 show that under reasonable assumptions (which, as we will see in Lemma 4.15, are satisfied by the class of meet-trees), the notions of a determined automorphism and an amalgamation base in the class \mathcal{K}_p^1 are essentially equivalent.

Proposition 4.2. *Suppose \mathcal{K} is a Fraïssé class with Fraïssé limit \mathbf{K} , and suppose p is a partial automorphism of \mathbf{K} such that for the structure $B \leq \mathbf{K}$ generated by $\text{dom}(p) \cup \text{range}(p)$, we have that (B, p) is an amalgamation base in \mathcal{K}_p^1 . Then p is determined.*

Proof. This is not important for our applications, so we only sketch the proof. We can do it by contraposition. If $f_1, f_2 \supseteq p$ are strict, not isomorphic over p , then there are some finite $f'_i \subseteq f_i$ ($i = 1, 2$) such that for no automorphism θ of \mathbf{K} fixing $\text{dom}(p)$ we have that $f'_1 \cup \theta f'_2 \theta^{-1}$

is a partial automorphism of \mathbf{K} . This easily implies that for the appropriate C_1, C_2 , we cannot amalgamate (C_1, f'_1) and (C_2, f'_2) over (B, p) . \square

Lemma 4.3. *Suppose \mathcal{K} is a Fraïssé class with Fraïssé limit \mathbf{K} . Consider the class $\overline{\mathcal{K}}$ of structures (possibly not finitely generated) whose age is contained in \mathcal{K} , and let $\overline{\mathcal{K}}^1$ be the class of $\overline{\mathcal{K}}$ -structures equipped with an automorphism.*

Then if $\overline{\mathcal{K}}^1$ has the AP, then for every determined partial automorphism p of \mathbf{K} , we have that if $B \subseteq \mathbf{K}$ is generated by $\text{dom}(p) \cup \text{range}(p)$, then (B, p) is an amalgamation base in \mathcal{K}_p^1 .

More generally, if for the unique strict extension $\hat{p} \supseteq p$, the structure $(\text{dom}(\hat{p}), \hat{p})$ is an amalgamation base in $\overline{\mathcal{K}}^1$, then (B, p) (defined as above) is an amalgamation base in \mathcal{K}_p^1 .

Proof. Consider two embeddings (in \mathcal{K}_p^1) $(B, p) \rightarrow (C_1, h_1), (C_2, h_2)$. Let $f_1, f_2 \in \text{Aut}(\mathbf{K})$ be strict extensions of h_1, h_2 (respectively). We may assume without loss of generality that the embeddings are simply inclusions. Then for $i = 1, 2$, we have a unique strict extension $f'_i \supseteq p$ such that $f'_i \subseteq f_i$, and we can take $B_i := \text{dom}(f'_i)$, and likewise, we can take $\overline{C}_i := \text{dom}(f_i)$. Since p is determined, we have an isomorphism $\theta_2: (B_1, f'_1) \rightarrow (B_2, f'_2)$ fixing $\text{dom}(p)$ (and hence B) pointwise. Now, for $i = 1, 2$, let $\overline{C}_i := \text{dom}(f_i)$. Write $(\overline{B}, f) := (B_1, f'_1)$ and write θ_1 for the inclusion mapping $\overline{B} \rightarrow \overline{C}_1$.

Note that clearly, $\overline{B}, \overline{C}_1, \overline{C}_2 \in \overline{\mathcal{K}}$, and θ_i yields an embedding $(\overline{B}, f) \rightarrow (\overline{C}_i, f_i)$ so, since (\overline{B}, f) is an amalgamation base in $\overline{\mathcal{K}}^1$, we have some $(D, h) \in \overline{\mathcal{K}}^1$ and embeddings $j_i: (\overline{C}_i, f_i) \rightarrow (D, h)$ (for $i = 1, 2$) such that $j_1 \circ \theta_1 = j_2 \circ \theta_2$. Now, put $D'_0 := j_1[\overline{C}_1] \cup j_2[\overline{C}_2]$ and let D' be the substructure of D generated by $D'_0 \cup h[D'_0]$. Put $h' := h \upharpoonright_{D'_0}$, $j'_i := j_i \upharpoonright_{\overline{C}_i}$.

Since $D \in \overline{\mathcal{K}}$, we have $D' \in \mathcal{K}$, so clearly $(D', h') \in \mathcal{K}_p^1$. Clearly, j'_i is an embedding of (C_i, h_i) into (D', h') , and since $j_1 \circ \theta_1 = j_2 \circ \theta_2$ and each θ_i fixes B pointwise, we also have

$$j'_1 \upharpoonright_B = j_1 \upharpoonright_B = (j_1 \circ \theta_1) \upharpoonright_B = (j_2 \circ \theta_2) \upharpoonright_B = j_2 \upharpoonright_B = j'_2 \upharpoonright_B.$$

Analogously, since $j_i \circ \theta_i \circ f \subseteq h \circ j_i \circ \theta_i$, we have $j'_i \upharpoonright_B \circ p \subseteq h' \circ j'_i \upharpoonright_B$, so (D', h') is an amalgam of (C_1, h_1) and (C_2, h_2) over (B, p) , which completes the proof. \square

Remark 4.4. The assumption that we have AP in $\overline{\mathcal{K}}^1$ (or at least that $(\text{dom}(\hat{p}), \hat{p})$ is an amalgamation base there) is necessary in Lemma 4.3. See Remark 4.19. \diamond

Remark 4.5. If $f \supseteq p$ is a strict extension, then f_+ , its restriction to the substructure generated by the positive f -orbits of elements of $\text{dom}(p)$, is a strictly positive extension. It is also easy to see that $\text{dom}(f) = \bigcup f^{-n} \text{dom}(f_+)$ (because it is a substructure of $\text{dom}(f)$ and it contains the f -orbits of elements of $\text{dom}(p)$). \diamond

We will be looking for determined partial automorphisms in order to apply Lemma 4.3. The following proposition shows that it is, in fact, enough to show positive determination.

Proposition 4.6. *If p is a positively determined finite partial automorphism, then it is determined.*

Proof. Fix two strict extensions $f, g \supseteq p$. Let f_+, g_+ be the respective positive parts. Fix an isomorphism $\theta_+: \text{dom}(f_+) \rightarrow \text{dom}(g_+)$ such that $g_+ = \theta_+ \circ f_+ \circ \theta_+^{-1}$ (which exists by the assumption). Put $\theta_n := g^{-n} \circ \theta_+ \circ f^n$.

Take any $a_+ \in \text{dom}(f_+)$. Then for any $k \geq 0$ we have:

$$f^k(a_+) = f_+^k(a_+) = \theta_+^{-1} \circ g_+^k \circ \theta_+(a_+) = \theta_+^{-1} \circ g^k \circ \theta_+(a_+),$$

and hence $g^{-k} \circ \theta_+ \circ f^k(a_+) = \theta_+(a)$. It follows that if $n \leq m$, then $\theta_n \subseteq \theta_m$. Indeed, if $a \in \text{dom}(\theta_n)$, then $f^n(a) \in \text{dom}(f_+)$, so $g^{n-m} \circ \theta_+ \circ f^{m-n}(f^n(a)) = \theta_+ \circ f^n(a)$, and thus

$$\theta_m(a) = g^{-n} \circ g^{n-m} \circ \theta_+ \circ f^{m-n}(f^n(a)) = g^{-n} \circ \theta_+ \circ f^n(a) = \theta_n(a).$$

It follows that $\theta := \bigcup_n \theta_n$ is a well-defined function. It is not hard to see that $\text{dom}(\theta) = \bigcup_n \text{dom}(\theta_n) = \text{dom}(f)$ and $\text{range}(\theta) = \bigcup_n (\text{range}(\theta_n)) = \text{dom}(g)$. Finally, since θ_+ is an isomorphism (between its domain and range), so is each θ_n , and hence also θ . \square

Definition 4.7. Let $p \subsetneq f$ be partial automorphisms of a structure M . We say that f is an *immediate extension* of p if $\text{dom}(f) \setminus \text{dom}(p)$ has only one element a , and moreover, $a \in \text{range}(p)$ and the p -orbit of a is the shortest non-cyclic orbit (i.e. no other non-cyclic orbit is shorter). \diamond

Remark 4.8. If p is a partial automorphism of a structure M , then p has a unique extension \bar{p} to a partial automorphism of M such that $\text{dom}(\bar{p})$ is the substructure of M generated by $\text{dom}(p)$. Furthermore, if $q \supseteq p$ is a partial automorphism extending p , then $\bar{q} \supseteq \bar{p}$. \diamond

Remark 4.9. If $f \supseteq p$ is a positively strict extension, then by straightforward induction, there is a sequence $(f_n)_n$ (with n ranging over ω or a finite ordinal) such that $f_0 = p$, for each $n > 0$, the extension $f_{n-1} \subseteq f_n$ is immediate, and $f = \bigcup_n \bar{f}_n$. \diamond

The following Lemma will be useful in showing that the determination of a partial automorphism.

Lemma 4.10. *Let M be a first order structure. Suppose p is a finite partial automorphism of M such that for every sequence $p = f_0 \subseteq f_1 \subseteq f_2 \subseteq \dots \subseteq f_n \subseteq f_{n+1}$ such that for all $i \leq n$, the extension $f_i \subseteq f_{i+1}$ is immediate, the following holds:*

- *If $g_n \supseteq f_n$ is positively strict, then there is an automorphism $\tau_n \in \text{Aut}(M)$ fixing $\text{dom}(f_n)$ pointwise, such that $\tau_n \circ f_{n+1} \circ \tau_n^{-1} \subseteq g_n$.*
- Then p is positively determined (and hence determined).*

Proof. Fix any two positively strict extensions $f, g \supseteq p$. For simplicity, suppose that f is infinite (the case when f is finite is analogous).

Let $(f_n)_n$ be a sequence as in Remark 4.9, so that $f_0 = p$ and $\bigcup_n \bar{f}_n = f$. We will recursively define a sequence θ_n of automorphisms of M such that:

- for all n , the domain of p is fixed pointwise by θ_n ,
- if $n \leq m$, then $\theta_n \upharpoonright_{\text{dom}(\bar{f}_n)} = \theta_m \upharpoonright_{\text{dom}(\bar{f}_n)}$,
- for all n we have $\theta_n \circ \bar{f}_n \circ \theta_n^{-1} \subseteq g$.

Then $\theta := \bigcup_n \theta_n \upharpoonright_{\text{dom}(\bar{f}_n)}$ will clearly be a well-defined embedding $\text{dom}(f) \rightarrow \text{dom}(g)$, fixing $\text{dom}(p)$, such that $\theta \circ f \circ \theta^{-1} \subseteq g$. It is easy to see that θ must be onto $\text{dom}(g)$ (because its image is a substructure which contains the positive g -orbits of elements of $\text{dom}(p)$), and hence $\theta \circ f \circ \theta^{-1} = g$. Since f, g are arbitrary, it will follow that p is positively determined, and hence (by Proposition 4.6) also determined.

It is clear that $\theta_0 = \text{id}_M$ satisfies the all the conditions listed above. Suppose we have $\theta_0, \dots, \theta_n$. Put $g_n := \theta_n^{-1} \circ g \circ \theta_n$. Then $g_n \supseteq p$ is positively strict (because $g \supseteq p$ is positively strict and θ_n fixes $\text{dom}(p)$ pointwise) and $g_n \supseteq f_n \supseteq p$, so g_n is a positively strict extension of

f_n . This allows us to take $\tau_n \in \text{Aut}(M)$ as in the hypothesis. Put $\theta_{n+1} := \theta_n \circ \tau_n$. Since τ_n fixes $\text{dom}(f_n)$, it follows that $\theta_{n+1} \upharpoonright_{\text{dom}(f_n)} = \theta_n \upharpoonright_{\text{dom}(f_n)}$, and hence also $\theta_{n+1} \upharpoonright_{\text{dom}(f_n)} = \theta_n \upharpoonright_{\text{dom}(f_n)}$.

Finally, since $\tau_n \circ f_{n+1} \circ \tau_n^{-1} \subseteq g_n = \theta_n^{-1} \circ g \circ \theta_n$, it easily follows that $\theta_{n+1} \circ f_{n+1} \circ \theta_{n+1}^{-1} \subseteq g$. \square

Determined partial automorphisms of trees

Now, we will proceed to show that determined partial automorphisms of finite trees are amalgamation bases (in the class of finite trees with partial automorphisms). To that end, we will show that the hypothesis of Lemma 4.3 is satisfied, i.e. that the class of meet-trees with an automorphism has the amalgamation property; we divide the proof into several steps.

Definition 4.11.

- We say that a meet-tree is *complete* if every chain has a least upper bound.
- Given a meet-tree T , its *completion* \hat{T} is the meet-tree consisting of cuts in T (i.e. the downwards closed chains, cf. Remark 2.22), ordered by inclusion, as in Remark 2.23. \diamond

(Note that for the particular case of linear orders, the definition of completeness given above is slightly more stringent than the usual one: it implies that there is a maximal element.)

Remark 4.12. If $(B, g), (C_1, f_1), (C_2, f_2)$ are trees with automorphisms, B is downwards closed in C_1 and C_2 , $B = C_1 \cap C_2$ and $f_1 \cap f_2 = g$, then $\hat{B} = \hat{C}_1 \cap \hat{C}_2$ and $\hat{g} = \hat{f}_1 \cap \hat{f}_2$, and $\hat{g}, \hat{f}_1, \hat{f}_2$ are automorphisms of $\hat{B}, \hat{C}_1, \hat{C}_2$, respectively (where $\hat{g}(\hat{b}) = g[\hat{b}]$ etc.). \diamond

Remark 4.13. If B_1, B_2 are sets, f_1 is a bijection on B_1 , f_2 is a bijection on B_2 , and f_1 and f_2 agree on $B = B_1 \cap B_2$, then $f_1 \cup f_2$ is a bijection on $B_1 \cup B_2$: because they agree, $f_1 \cup f_2$ is a well-defined function, and it is easy to check that $f_1^{-1} \cup f_2^{-1}$ is its inverse. \diamond

Lemma 4.14. *Suppose $(B, g), (C_1, f_1), (C_2, f_2)$ are trees with automorphisms such that $B = C_1 \cap C_2$ and $g = f_1 \cap f_2$. Then there are $B', C'_1, C'_2, g', f'_1, f'_2$ such that:*

- $B' \subseteq C'_1, C'_2$ and is downwards closed in both,
- g' is automorphism of B and f'_i is an automorphism of C'_i for $i = 1, 2$,
- $(B, g) \subseteq (B', g')$ and $(C_i, f_i) \subseteq (C'_i, f'_i)$ for $i = 1, 2$.

Proof. Let B'_i be the downwards closure of B in C_i for $i = 1, 2$. Put $B' := B'_1 \cup B'_2$, $g' := f_1 \upharpoonright_{B'_1} \cup f_2 \upharpoonright_{B'_2}$, $C'_i := C_i \cup B'$ and $f'_i := f_i \cup g'$ for $i = 1, 2$.

We put on each B'_i the meet-tree structure inherited from C_i . Recall from Remark 2.23 that we have a canonical ordering on the cuts in B . Now, take some c_1, c_2 from $B'_1 \setminus B$ and $B'_2 \setminus B$, respectively, and denote by q_1, q_2 their cuts in B , and take any $b_1, b_2 \in B$ such that $b_i > q_i$. Then:

- if $q_1 \leq q_2$, we declare $c_1 < c_2$ and $c_1 \wedge c_2 = c_1$,
- if $q_1 > q_2$, we declare $c_1 > c_2$ and $c_1 \wedge c_2 = c_2$,
- if q_1 and q_2 are incomparable, we declare that c_1 and c_2 are incomparable and $c_1 \wedge c_2 = b_1 \wedge b_2$.

The fact that this definition of $<$ defines a semilinear partial order is left as an exercise. Let us only show that the \wedge is correct and well-defined in the last case. We need to show that given $c \in B'$, we have that $c \leq c_1, c_2$ if and only if $c \leq b_1 \wedge b_2$ (this immediately implies that $c_1 \wedge c_2$ does not depend on the choice of b_1 and b_2). The fact that the left-hand side implies the right-hand side is trivial, since $b_i \geq c_i$. For the converse, suppose $c \leq b_1 \wedge b_2$. Note that c_1 and c are comparable, so it is enough to show that $c \not\leq c_1$. But $c \leq b_2$, and we cannot have $c_1 \leq b_2$ (because c_1 and $c_2 \leq b_2$ are incomparable).

Now, we need to check that g' is an automorphism of B' . Now, it is clear that $f_i \upharpoonright_{B'_i}$ is an automorphism of B'_i for each i , and by Remark 4.13, g' is a bijection, so we only need to check that it preserves the ordering and the meets. But this is easy to see as a consequence of the observation that g' preserves $B, B'_1 \setminus B, B'_2 \setminus B$, and the cuts of all elements of B' .

Finally, again by Remark 4.13, each f'_i is a bijection on C'_i , and the fact that it is an automorphism follows from Lemma 3.5 (with $f = g'$ and $g = f_i$). \square

Lemma 4.15. *The class of meet-trees with automorphisms has the amalgamation property.*

Proof. Suppose $(B, g), (C_1, f_1), (C_2, f_2)$ are trees with automorphisms, with fixed embeddings of (B, g) into each (C_i, f_i) . We need to find a tree with automorphism (D, h) such that (C_i, f_i) embed in (D, h) for $i = 1, 2$ in such a way that the two resulting embeddings of (B, g) into (D, h) coincide.

We may assume for simplicity that $(B, g) \subseteq (C_i, f_i)$ for $i = 1, 2$, while $C_1 \cap C_2 = B$ and $f_1 \cap f_2 = g$. We will find a (D, h) such that $(C_i, f_i) \subseteq (D, h)$ (which will immediately imply that the two embeddings of (B, g) coincide).

By Lemma 4.14, we may assume without loss of generality that B is downwards closed in C_1, C_2 . Then by Remark 4.12, we may also assume that B, C_1, C_2 are complete. We put $D := C_1 \cup C_2$ and $h := f_1 \cup f_2$. We need to describe the meet-tree structure on D and to show that h is an automorphism of this structure. Now, for each $c \in C_1 \cup C_2$, put $b_c := \sup\{b \in B \mid b \leq c\}$ (this is well-defined by completeness of B).

On D , we define the structure in the following way:

- on each C_i , the structure is simply the original structure,
- given $c_1 \in C_1 \setminus B$ and $c_2 \in C_2 \setminus B$, we declare c_1 and c_2 to be incomparable and put $c_1 \wedge c_2 := b_{c_1} \wedge b_{c_2}$ (where the meet on the right-hand side is in the sense of B).

Note that this gives a meet-tree structure: indeed, given any $d_0 \in D$, the interval $(\infty, d_0]$ in D is contained in C_1 or C_2 (depending on whether $d_0 \in C_1$ or $d_0 \in C_2$), so semilinearity follows from semilinearity of C_1 and C_2 . By the same token, given $c_1 \in C_1$ and $c_2 \in C_2$, the intersection $(-\infty, c_1] \cap (-\infty, c_2]$ is contained in B , so it is contained in $(-\infty, b_{c_1}] \cap (-\infty, b_{c_2}] = (-\infty, b_{c_1} \wedge b_{c_2}]$.

Furthermore, h is an automorphism of D : by Remark 4.13, it is a bijection, and it clearly preserves meets and inequalities within each C_i . Now, given $c_1 \in C_1 \setminus C_2$ and $c_2 \in C_2 \setminus C_1$, the two are incomparable, and since $h(c_1) = f_1(c_1) \in C_1 \setminus C_2$ and $h(c_2) = f_2(c_2) \in C_2 \setminus C_1$ (because C_1, C_2 are clearly h -invariant), by the same token, $h(c_1)$ and $h(c_2)$ are incomparable. It is also not hard to see that $b_{f_i(c_i)} = f_i(b_{c_i})$, which implies that h preserves meets.

Finally, clearly $D \supseteq C_1, C_2$ and $h \supseteq f_1, f_2$, so (D, h) is as desired. \square

One could ask whether the analogue of Lemma 4.15 for trees of bounded arity is true. Unfortunately, this is not the case, which the following proposition shows.

Proposition 4.16. *For every integer $k > 1$, the class of trees of arity at most k with an automorphism does not have the amalgamation property.*

Proof. Fix k and let $B = \{b\}$, $C_1 = B \cup \{c_1^1, c_1^2, \dots, c_1^k\}$, $C_2 = B \cup \{c_2\}$ be meet-trees such that $c_2 > b$ and for all $i \neq j$ we have $c_1^i, c_1^j > c_1^i \wedge c_1^j = b$. Define f on B as $f(b) = b$, $f_1 \supseteq f$ on C_1 as $f_1(c_1^i) = c_1^{i+1}$, for $i < k$, $f_1(c_1^k) = c_1^1$, and finally define $f_2 \supseteq f$ on C_2 as the identity map.

We claim that (C_1, f_1) and (C_2, f_2) do not amalgamate over (B, f) . Indeed, suppose we do have an amalgam (D, g) , with D being a tree of arity at most k . It follows that for some i we have $c_1^i \wedge c_2 > b$. We may assume without loss of generality that $i = 1$. But then by Fact 2.18, since $b = c_1^1 \wedge c_1^2$, it follows that $c_1^2 \wedge c_2 = b$. On the other hand, since g is an automorphism of D , we have $b = g(b) < g(c_1^1 \wedge c_2) = g(c_1^1) \wedge g(c_2) = c_1^2 \wedge c_2$, a contradiction. \square

In the opposite direction to Proposition 4.16, the following Remark describes a possible class of amalgamation classes for the case of bounded arity.

Remark 4.17. We suspect that the following is true (but, as we do not use it, we will did not check it very carefully): if $k > 1$ and (B, g) is a meet-tree of arity at most k with an automorphism, such that every g -periodic element of b has maximal rank (i.e. there are b_1, \dots, b_k which are pairwise incomparable, such that $b_i \wedge b_j = b$ for all $i \neq j$), then (B, g) is an amalgamation base (in the class of all meet-trees of arity at most k with an automorphism). \diamond

The following corollary summarises the results of this section. In the rest of the paper, we will show (using Lemma 4.10) that determined finite partial automorphisms of \mathbf{T} are cofinal, which will imply that \mathbf{T} has a generic automorphism.

Corollary 4.18. *If p is a determined finite partial automorphism of \mathbf{T} and $B \subseteq \mathbf{T}$ is generated by $\text{dom}(p) \cup \text{range}(p)$, then (B, p) is an amalgamation base in the class of finite trees with partial automorphisms.*

Proof. Note that if the age of a structure consists of finite meet-trees, then it is a meet-tree itself, and conversely, the age of any meet-tree consists of finite meet-trees. The corollary follows immediately by Lemma 4.15 and Lemma 4.3. \square

Remark 4.19. Note that the construction from Proposition 4.16 actually shows that the analogue of Corollary 4.18 for \mathbf{T}_k fails (when $k > 1$): an automorphism of \mathbf{T}_k which is just a single fixed point is trivially determined, and the example shows that it does not yield an amalgamation base. \diamond

Remark 4.20. One can show that the class of linear orders with an automorphism does have AP (essentially, arguing as in the proof of Lemma 4.14), so the analogue of Corollary 4.18 for \mathbf{T}_1 holds. This can be used (along with a variant of Theorem 6.6) to recover the fact that $(\mathbf{Q}, <)$ has a generic automorphism, see Remark 6.13. \diamond

5 Orbits in meet-trees

Having Corollary 4.18, to show CAP for the class of finite automorphisms of meet-trees, it is enough to show that determined partial automorphisms are cofinal. Before we can do that, we need to understand the orbits of partial automorphisms in trees and their extensions.

Definition 5.1. A (finite) *partial orbit* (in a meet-tree) is a finite sequence $\eta = (\eta_0, \eta_1, \dots, \eta_n)$ of elements of the tree such that there is a partial automorphism p such that for each $i = 0, \dots, n-1$ we have $p(\eta_i) = \eta_{i+1}$. Infinite partial orbits are defined analogously. \diamond

Until the end of the paper, we will use the convention that lowercase Greek letters represent (usually finite) partial orbits in meet-trees, and each such orbit η is enumerated as η_0, \dots, η_n , as in Definition 5.1.

Remark 5.2. Since every finite tree can be embedded in \mathbf{T} (which is ultrahomogeneous), every finite partial meet-tree automorphism can be extended to a total meet-tree automorphism (possibly after enlarging the tree). In particular, every finite partial orbit is contained in the orbit of some total automorphism.

(Likewise, every finite automorphism and finite partial orbit in a meet-tree of arity at most k can be extended to an automorphism or a full orbit of a meet-tree of arity at most k .) \diamond

Definition 5.3. Let $\eta = (\eta_0, \eta_1, \dots, \eta_n)$ be a partial orbit, while k is a positive integer.

Then:

- We say that η is a *k-cycle* (or *k-cyclic*) if k is minimal such that for some $\eta_k = \eta_0$.
- We say that η is an *ascending k-spiral* if k is minimal such that $\eta_k > \eta_0$. Likewise, we say that it is a *descending k-spiral* if k is minimal such that $\eta_k < \eta_0$.
- We say that η is an *ascending k-comb* if it is not a spiral and k is minimal such that $\eta_{2k} \wedge \eta_k > \eta_k \wedge \eta_0$. Likewise, we say that it is a *descending k-comb* if it is not a spiral and k is minimal such that $\eta_{2k} \wedge \eta_k < \eta_k \wedge \eta_0$.
- Otherwise, if η is not a cycle, spiral, nor a comb, we say that η is a *quasi-cycle* (or *quasi-cyclic*).

If η is a k -cycle, we say that k is the *period* of η ; if η is a k -spiral or a k -comb, we say that k is its *spiral length*.

We define the *length* of η as the size of $\{\eta_0, \dots, \eta_n\}$.

(See Figures 1 and 2 for some examples of orbits of various types.) \diamond

Remark 5.4. If $\eta \subseteq \bar{\eta}$ are finite partial orbits, then if η is a k -cycle, an ascending/descending k -spiral or k -comb, then so is $\bar{\eta}$ (respectively). On the other hand, if η is a quasi-cycle, then it does not say much about the kind of orbit $\bar{\eta}$ can be (see Remark 5.18). \diamond

Remark 5.5. If $\eta = (\eta_0, \dots, \eta_n)$ is a partial orbit, then so is $\eta^{-1} = (\eta_n, \eta_{n-1}, \dots, \eta_0)$. If we think of η as a sequence of steps in a process, then η^{-1} can be thought of as “time-reversal” of this process. \diamond

The following proposition describes the spiral behaviour of orbits.

Proposition 5.6. *Suppose that $\eta = (\eta_0, \dots, \eta_n)$ is k -spiral. Then for $i, j, l \in [0, n]$:*

- (1) η_i is comparable to η_j if and only if $i \equiv j \pmod{k}$,
- (2) if $i \equiv j \not\equiv l \pmod{k}$, then $\eta_i \wedge \eta_l = \eta_j \wedge \eta_l$.

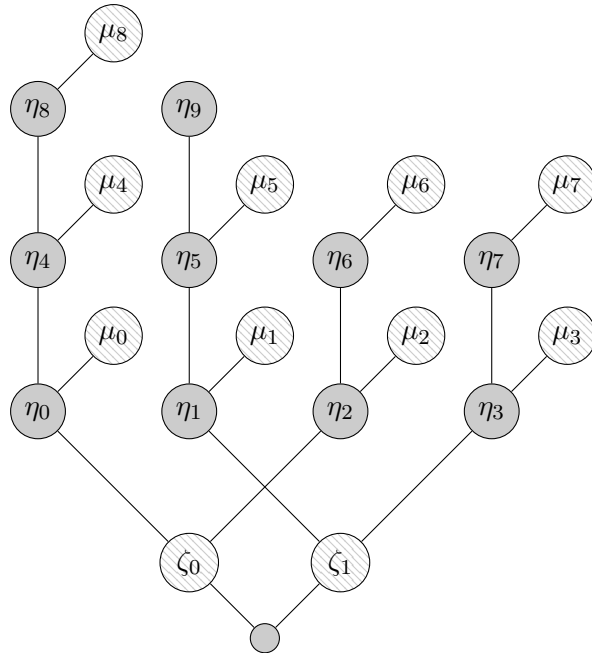


Figure 1: In the above figure, the ζ -orbit is a 2-cycle, the η -orbit is an ascending 4-spiral, while the μ -orbit is an ascending 4-comb. The root is a fixed point.

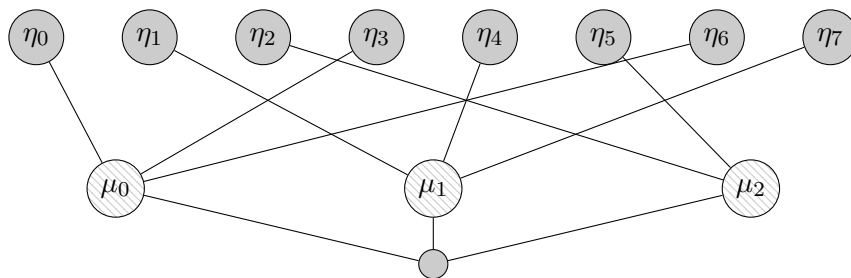


Figure 2: In the above figure, the η -orbit is a quasi-cycle, the μ -orbit is a 3-cycle, while the root is a fixed point.

Proof. The case of $k = 1$ is straightforward, so we may assume that $k > 1$. We may assume without loss of generality that η is ascending (replacing it with η^{-1} if necessary), and we may also assume that $n \geq jk$ (extending η if necessary).

(1): It is clear (by straightforward induction) that if $i \equiv j \pmod{k}$, then η_i and η_j are comparable. For the converse, we may assume without loss of generality that $i < j$, and in fact that $i = 0$. Since η is an ascending k -spiral, $\eta_0 < \eta_{jk}$. This clearly implies that $\eta_j \not\leq \eta_0$ (otherwise, we would have $\eta_{jk} \leq \eta_0$), so by comparability, $\eta_j > \eta_0$. Let m be maximal such that $km \leq j$. Since $j > 0$ and k is minimal such that $\eta_k > \eta_0$, we have $m > 0$. Now, it is easy to see that we have $\eta_{km}, \eta_j < \eta_{jk}$, so η_{km}, η_j are comparable. It follows that η_0, η_{j-km} are comparable. But $0 \leq j - km < k$, so by the choice of k , it follows that $j - km = 0$, so k divides j .

(2): By (1), η_i and η_j are comparable, so without loss of generality we can assume $\eta_i \leq \eta_j$. It follows that $\eta_i \wedge \eta_j = \eta_i$. Also by (1), η_i, η_l are not comparable, so $\eta_i \wedge \eta_l < \eta_i$. Together, we have that $\eta_i \wedge \eta_j > \eta_i \wedge \eta_l$, so by Fact 2.18, $\eta_i \wedge \eta_l = \eta_j \wedge \eta_l$. \square

Proposition 5.7. *Let $\eta = (\eta_0, \dots, \eta_m)$ be a partial orbit. Take any positive $k \leq n/2$ and integers $i_1 \equiv i_2 \equiv j_1 \equiv j_2 \pmod{k}$ such that $i_1 < i_2$ and $j_1 < j_2$.*

If $\eta_0 \wedge \eta_k < \eta_k \wedge \eta_{2k}$ (in particular, if η is an ascending k -spiral or k -comb), then:

- $\eta_{i_1} \wedge \eta_{i_2} = \eta_{i_1} \wedge \eta_{i_1+k}$;
- $\text{otp}(\eta_{i_1} \wedge \eta_{i_2}, \eta_{j_1} \wedge \eta_{j_2}) = \text{otp}(i_1, j_1)$.

Likewise, if $\eta_0 \wedge \eta_k > \eta_k \wedge \eta_{2k}$ (in particular, if η is a descending k -spiral or k -comb), then:

- $\eta_{i_1} \wedge \eta_{i_2} = \eta_{i_2-k} \wedge \eta_{i_2}$;
- $\text{otp}(\eta_{i_1} \wedge \eta_{i_2}, \eta_{j_1} \wedge \eta_{j_2}) = \text{otp}(j_2, i_2)$.

Proof. We will consider the ascending case (the descending case is completely analogous, and in fact it follows by considering the “time reversal”). The first bullet follows by induction with respect to $i_2 - i_1$; if $i_2 - i_1 = k$, then the conclusion is trivial. Otherwise, if $i_2 > i_1 + k$, by induction hypothesis, we have that $\eta_{i_1+k} \wedge \eta_{i_2} = \eta_{i_1+k} \wedge \eta_{i_1+2k}$. Since $\eta_0 \wedge \eta_k < \eta_k \wedge \eta_{2k}$, also $\eta_{i_1} \wedge \eta_{i_1+k} < \eta_{i_1+k} \wedge \eta_{i_1+2k}$, whence $\eta_{i_1} \wedge \eta_{i_1+k} < \eta_{i_1+k} \wedge \eta_{i_2}$, and the conclusion follows by Fact 2.18.

For the second bullet, by the first bullet, we may assume without loss of generality that $i_2 = i_1 + k$ and $j_2 = j_1 + k$. We may also assume without loss of generality that $i_1 < j_1$. Straightforward induction shows that $\eta_{i_1} \wedge \eta_{i_1+k} < \eta_{j_1} \wedge \eta_{j_1+k}$, which finishes the proof. \square

Corollary 5.8. *Fix a meet-tree T , and arbitrary $a_1, a_2 \in T$, as well as $B \subseteq T$. Then if $a_1 \wedge a_2 > \max_{b \in B} b \wedge a_1$, then $\text{qftp}(a_1/B) = \text{qftp}(a_2/B)$ (in particular, if $B = M$, then $\text{tp}(a_1/B) = \text{tp}(a_2/B)$).*

Proof. Note that the assumption immediately implies that $a_1, a_2 \neq \max_{b \in B} b \wedge a_1$, so by Fact 2.24 it is enough to show that $\max_{b \in B} b \wedge a_2 = \max_{b \in B} b \wedge a_1$. But the assumption tells us immediately that $b \wedge a_1 < a_1 \wedge a_2$, so by Fact 2.18, $b \wedge a_1 = b \wedge a_2$, which completes the proof. \square

Remark 5.9. If p is a partial automorphism of a first-order structure C and $a, b \in C$, then $p \cup \{(a, b)\}$ is a partial automorphism of C if and only if $\text{qftp}(b/\text{range}(p)) = p(\text{qftp}(a/\text{dom}(p)))$. \diamond

Definition 5.10. Given a finite partial orbit $\eta = (\eta_0, \dots, \eta_n)$ ($n > 0$), its *pseudo-period* is the smallest $u > 0$ such that $\eta_0 \wedge \eta_u = \max_{0 < i \leq n} \eta_0 \wedge \eta_i$. \diamond

Remark 5.11. Note that if η is a cycle, then its period is also its pseudo-period. \diamond

Proposition 5.12. *The pseudo-period is invariant under time reversal.*

More precisely, if $\eta = (\eta_0, \dots, \eta_n)$ is a partial orbit of pseudo-period u , then u is also the pseudo-period of $\eta^{-1} = (\eta_n, \eta_{n-1}, \dots, \eta_1, \eta_0)$, i.e. it is the smallest $m_0 > 0$ such that $\eta_{n-m_0} \wedge \eta_n = \max_{0 < m \leq n} \eta_n \wedge \eta_{n-m}$.

Proof. Let u' be the pseudo-period of η^{-1} . By symmetry, it is enough to show that $u' \geq u$. This is equivalent to saying that for all positive $m < u$, we have $\eta_n \wedge \eta_{n-m} < \eta_n \wedge \eta_{n-u}$. Note that for such m we have $0 < u-m < u$, so $\eta_0 \wedge \eta_{u-m} < \eta_0 \wedge \eta_u$. Now, let f be an automorphism of a meet-tree such that for $f(\eta_i) = \eta_{i+1}$ for $i = 0, \dots, n-1$. By applying f^{n-u} , we obtain that $\eta_{n-u} \wedge \eta_{n-m} < \eta_{n-u} \wedge \eta_n$. By Fact 2.18, it follows that $\eta_{n-u} \wedge \eta_{n-m} = \eta_n \wedge \eta_{n-m}$ and so we are done. \square

Proposition 5.13. *Suppose $\eta = (\eta_0, \dots, \eta_n)$ is a quasi-cycle of pseudo-period u . Then either $n \geq 2u$ or η can be extended to $\bar{\eta} = (\eta_0, \dots, \eta_n, \eta'_{n+1}, \dots, \eta'_{2u})$, a partial orbit such that $\eta'_{2u} = \eta_0$ (so $\bar{\eta}$ is a $2u$ -cycle).*

As a consequence, η is the orbit of a finite partial automorphism p such that $\eta_0 \wedge \eta_u$ is a fixed point of p^u (in particular, $p^u(\eta_0 \wedge \eta_u) \downarrow$).

Proof. Let f be a meet-tree automorphism, one of whose orbits includes η . For each $i > n$, put $\eta_i := f^i(\eta_0)$. We may assume without loss of generality that $n+1 \leq 2u$ (otherwise, the conclusion is trivial).

Claim. $(\eta_0, \dots, \eta_{2u-1})$ is a quasi-cycle.

Proof. We will prove by induction that (η_0, \dots, η_m) is a quasi-cycle, where $m < 2u$. If $m \leq n$, then there is nothing to prove. Suppose now that $n \leq m < 2u-1$ and we know that (η_0, \dots, η_m) is a quasi-cycle, and we show that so is $(\eta_0, \dots, \eta_{m+1})$.

First, we show that $\eta_{m+1} \not\leq \eta_0$ — the argument for $\eta_0 \not\leq \eta_{m+1}$ is symmetric, in light of Proposition 5.12, and the two together show that we have neither a cycle nor a spiral. Note that since $m+1 < 2u$, we have $m+1-u < u$. It follows immediately that $\eta_0 \wedge \eta_u > \eta_0 \wedge \eta_{m+1-u}$, and by Proposition 5.12, it also follows that $\eta_{m+1} \wedge \eta_{m+1-u} > \eta_{m+1} \wedge \eta_u$. This clearly implies that either $\eta_0 \wedge \eta_{m+1-u} \not\leq \eta_{m+1} \wedge \eta_{m+1-u}$ or $\eta_{m+1} \wedge \eta_u \not\leq \eta_0 \wedge \eta_u$.

Suppose first that $\eta_0 \wedge \eta_{m+1-u} \not\leq \eta_{m+1} \wedge \eta_{m+1-u}$. By semilinearity, it follows that $\eta_0 \wedge \eta_{m+1-u} < \eta_{m+1} \wedge \eta_{m+1-u}$, so by Fact 2.18, $\eta_0 \wedge \eta_{m+1-u} = \eta_0 \wedge \eta_{m+1}$. But note that since η is a quasi-cycle, $\eta_0 \wedge \eta_{m+1-u} < \eta_0$, so it follows that $\eta_0 \not\leq \eta_{m+1}$. Otherwise, if $\eta_{m+1} \wedge \eta_u \not\leq \eta_0 \wedge \eta_u$, then analogously $\eta_{m+1} \wedge \eta_u = \eta_0 \wedge \eta_{m+1}$, and since $\eta_m > \eta_m \wedge \eta_{u-1}$, also $\eta_{m+1} > \eta_{m+1} \wedge \eta_u = \eta_0 \wedge \eta_{m+1}$, so $\eta_{m+1} \not\leq \eta_0$.

Now, if, in addition, we have that m is odd, i.e. $m+1 = 2c$ for some $c < u$, then we need to show that $\eta_0 \wedge \eta_c = \eta_c \wedge \eta_{2c}$. Since $c < u \leq m = 2c-1$, we have $2c > u$, so $c > u-c$, whence $m = 2c-1 \geq 2(u-c)$. Thus, by the induction hypothesis, $\eta_0 \wedge \eta_{u-c} = \eta_{u-c} \wedge \eta_{2u-2c}$. On the other hand, $0 < u-c < u$, so $\eta_0 \wedge \eta_{u-c} < \eta_0 \wedge \eta_u$, whence $\eta_0 \wedge \eta_{u-c} = \eta_u \wedge \eta_{u-c}$, so $\eta_u \wedge \eta_{u-c} = \eta_{u-c} \wedge \eta_{2u-2c}$. By applying p^{2c-u} , we obtain $\eta_{2c} \wedge \eta_c = \eta_c \wedge \eta_u$. But since $c < u$, we have also $\eta_0 \wedge \eta_c < \eta_0 \wedge \eta_u$, so by another application of Fact 2.18, $\eta_0 \wedge \eta_c = \eta_u \wedge \eta_c$, which completes the proof. \square (claim)

We will show that $\bar{\eta} = (\eta_0, \dots, \eta_{2u-1}, \eta_0)$ is a partial orbit (which is clearly sufficient).

By Remark 5.9, to complete the proof, it is enough to show that $\text{tp}(\eta_0/\eta_1, \dots, \eta_{2u-1}) = \text{tp}(\eta_{2u}/\eta_1, \dots, \eta_{2u-1})$. Since (η_0, \dots, η_m) is a quasi-cycle, it follows that $\eta_0, \eta_u > \eta_0 \wedge \eta_u$, and hence also $\eta_u, \eta_{2u} > \eta_u \wedge \eta_{2u}$. We will show that for every $0 < i, j \leq m$, we have that $b := \eta_i \wedge \eta_j < \eta_0 \wedge \eta_u$ if and only if $b < \eta_u$, if and only if $b < \eta_u \wedge \eta_{2u}$. Since η is a quasi-cycle, it easily follows that we cannot have $b > \eta_u$, so by Fact 2.24, it will follow that the types are equal as required.

It is clear that if $b < \eta_0 \wedge \eta_u$ or $b < \eta_u \wedge \eta_{2u}$, then $b < \eta_u$, so suppose that $b < \eta_u$, and let us show that $b < \eta_u \wedge \eta_{2u}$ and $b < \eta_0 \wedge \eta_u$. Note that if $b = \eta_i \wedge \eta_j$, then by Fact 2.18, it follows that either $b = \eta_u \wedge \eta_i$ or $b = \eta_u \wedge \eta_j$. Without loss of generality we may assume the former.

Now, since $b \neq u$, we have $i \neq u$. Let us consider the case when $i > u$ (the other case is analogous). Write k for $i - u$, so that $i = u + k$. Clearly, $\eta_0 \wedge \eta_u > \eta_0 \wedge \eta_k$, so $\eta_u \wedge \eta_{2u} > \eta_{u+k} \wedge \eta_u$. Furthermore, $0 < k < u$, so $\eta_0 \wedge \eta_u > \eta_0 \wedge \eta_{u-k}$, so $\eta_0 \wedge \eta_{u-k} = \eta_u \wedge \eta_{u-k}$. whence $\eta_k \wedge \eta_u = \eta_{u+k} \wedge \eta_u$. But $\eta_0 \wedge \eta_u > \eta_0 \wedge \eta_k$, so $\eta_0 \wedge \eta_k = \eta_u \wedge \eta_k = \eta_{u+k} \wedge \eta_u$ and we are done. \square

Proposition 5.14. *Let $\eta = (\eta_0, \dots, \eta_n)$ be a quasi-cycle with pseudo-period u .*

Then if $i, j, k \in \{0, \dots, n\}$ satisfy $i \equiv j \pmod{u}$ and $k \neq i, j$, then $\eta_i \wedge \eta_k = \eta_j \wedge \eta_k$.

Proof. By Proposition 5.12, we may assume that $j < k$. Indeed, otherwise, if $j > k$, then we can simply consider η^{-1} — this preserves u , but reverses the order of η_j and η_k in the orbit. Further, we may also assume without loss of generality that $j < i$ (otherwise, if $i < j < k$, we can just swap i and j ; the case of $i = j$ is trivial). Moreover, it is enough to consider the case when $j = 0$ (truncating η if necessary). This leaves us with some $i, k > 0$ such that $i \neq k$ and u divides i , and we need to show that $\eta_0 \wedge \eta_k = \eta_i \wedge \eta_k$.

By Proposition 5.13, we have a finite automorphism p such that η is an orbit of p and $p^u(\eta_0 \wedge \eta_u) = \eta_0 \wedge \eta_u$. Let $f \supseteq p$ be a total meet-tree automorphism.

Claim. *If $0 < m \leq n$, then $\eta_0 \wedge \eta_m = \eta_0 \wedge \eta_u$ if and only if u divides m .*

Proof. Suppose first that u divides m , so $m = lu$. The proof is by induction with respect to l . Since η is a quasi-cycle, we have that $\eta_{lu-u} \wedge \eta_{lu} = \eta_{lu-2u} \wedge \eta_{lu-u} = \dots = \eta_0 \wedge \eta_u$, so in particular, $\eta_0 \wedge \eta_u = \eta_{lu-u} \wedge \eta_{lu}$. By the induction hypothesis, $\eta_0 \wedge \eta_u = \eta_0 \wedge \eta_{lu-u}$, so $\eta_0 \wedge \eta_{lu-u} = \eta_{lu-u} \wedge \eta_{lu}$. By Fact 2.18, it follows that $\eta_0 \wedge \eta_{lu} \geq \eta_0 \wedge \eta_{lu-u} = \eta_0 \wedge \eta_u$. Since by the definition of u , $\eta_0 \wedge \eta_u \geq \eta_0 \wedge \eta_{lu}$, this completes the proof.

Now, suppose $\eta_0 \wedge \eta_m = \eta_0 \wedge \eta_u$. Let l be maximal such that $lu \leq m$. Then, by the preceding paragraph, $\eta_0 \wedge \eta_{lu} = \eta_0 \wedge \eta_u$, so by Fact 2.18, $\eta_{lu} \wedge \eta_m \geq \eta_0 \wedge \eta_u$. Now, $\eta_0 \wedge \eta_u$ is a fixed point of f^u , and hence also of p^{-lu} . By applying f^{-lu} to the last inequality, we obtain $f^{-lu}(\eta_{lu} \wedge \eta_m) = \eta_0 \wedge \eta_{m-lu} \geq \eta_0 \wedge \eta_u$. But $m - lu < u$, so by minimality of u , we have that $m - lu = 0$, so $m = lu$. \square (claim)

Now, by Claim, we have $\eta_0 \wedge \eta_i = \eta_0 \wedge \eta_u \geq \eta_0 \wedge \eta_k$. If the inequality is strict, by Fact 2.18, $\eta_0 \wedge \eta_k = \eta_i \wedge \eta_k$ and we are done. Otherwise, by Claim, u divides k and $\eta_0 \wedge \eta_k = \eta_0 \wedge \eta_u$, so we need to show that $\eta_i \wedge \eta_k = \eta_0 \wedge \eta_u$.

We have that u divides $k - i$ so by Claim, we have $\eta_0 \wedge \eta_{|k-i|} = \eta_0 \wedge \eta_u$. Since $\eta_0 \wedge \eta_u$ is a fixed point of p^u and u divides i, k (and hence also $\min(i, k)$), $\eta_0 \wedge \eta_u$ is also a fixed point of $p^{\min(i, k)}$, so we have that

$$\eta_i \wedge \eta_k = \eta_{\min(i, k)} \wedge \eta_{\min(i, k) + |i - k|} = p^{\min(i, k)}(\eta_0 \wedge \eta_u) = \eta_0 \wedge \eta_u,$$

so we are done. \square

Corollary 5.15. *If $\eta = (\eta_0, \dots, \eta_m)$ is a quasi-cycle with pseudo-period u , and $N > n$ is minimal such that u divides N , then every extension $\bar{\eta} = (\eta_0, \dots, \eta_{N-1})$ is also a quasi-cycle with pseudo-period u .*

Proof. The case of $N = 2u$ is essentially contained in Proposition 5.13, so we may assume that $n \geq 2u$.

First, we will show that the pseudo-period is indeed u . Take any $m \in (n, N)$. Then u does not divide m and $2u \leq n < m \leq n + u$, so (η_u, \dots, η_m) is a quasi-cycle with pseudo-period u ; thus, Proposition 5.14 implies that $\eta_u \wedge \eta_m = \eta_u \wedge \eta_{m'}$, where $u < m' < 2u$, so $\eta_u \wedge \eta_m = \eta_u \wedge \eta_{m'} < \eta_u \wedge \eta_{2u}$. But by the same token, $\eta_u \wedge \eta_{2u} = \eta_0 \wedge \eta_u$, so $\eta_u \wedge \eta_m < \eta_0 \wedge \eta_u$, and so, by Fact 2.18, $\eta_u \wedge \eta_m = \eta_0 \wedge \eta_m < \eta_0 \wedge \eta_u$.

Now, we will show that (η_0, \dots, η_m) is a quasi-cycle by induction with respect to $m \geq n$. The case of $m = n$ is clear. Suppose now that $n < m < N - 1$ and (η_0, \dots, η_m) is a quasi-cycle (with pseudo-period u). We need to show that so is $(\eta_0, \dots, \eta_{m+1})$.

It easily follows from the assumption that u does not divide $m + 1$, so we have some positive $k \in (u, 2u)$ such that $k \equiv m + 1 \pmod{u}$. By Proposition 5.14, it follows that $\eta_u \wedge \eta_{m+1} = \eta_u \wedge \eta_k < \eta_u \wedge \eta_{2u} = \eta_0 \wedge \eta_u < \eta_0$. It follows that $\eta_{m+1} \not\geq \eta_0$. The not-inequality $\eta_{m+1} \not\leq \eta_0$ can be proven similarly, or follows immediately from the above argument by considering the partial orbit $(\eta_{m+1}, \eta_m, \dots, \eta_0)$ (and applying Proposition 5.12). In conclusion, $(\eta_0, \dots, \eta_{m+1})$ is not a cycle, and not a spiral. We need to prove that it is not a comb.

Suppose now that m is odd, so $m + 1 = 2c$. We need to show that $\eta_0 \wedge \eta_c = \eta_c \wedge \eta_{2c}$. Let $k < u$ be such that $k \equiv c \pmod{u}$ (since u does not divide $2c$, we have $k > 0$, and u divides none of $k, 2k, c, 2c$). Then clearly $c, k \neq 0$, $c \neq 2c$ and $k \neq 2k$, so (since $(\eta_1, \dots, \eta_{m+1})$ is a quasi-cycle with pseudo-period u), by Proposition 5.14, $\eta_c \wedge \eta_{2c} = \eta_k \wedge \eta_{2k}$. Again by Proposition 5.14, $\eta_0 \wedge \eta_c = \eta_0 \wedge \eta_k$. Since $k < u$, $2k < 2u \leq n$, so $\eta_0 \wedge \eta_k = \eta_k \wedge \eta_{2k}$, whence $\eta_0 \wedge \eta_c = \eta_c \wedge \eta_{2c}$. \square

Corollary 5.16. *If $\eta = (\eta_0, \dots, \eta_n)$ is a quasi-cycle, then it can be extended to a cycle.*

Proof. Let u be the pseudo-period of η . The case of $n < 2u$ is Proposition 5.13, so we may assume that $n \geq 2u$. Let N be as in Corollary 5.15, and let (η_0, \dots, η_N) be an arbitrary extension of η . We will show that $(\eta_0, \dots, \eta_{N-1}, \eta_0)$ is a partial orbit. Write B for the subtree generated by $\{\eta_1, \dots, \eta_{N-1}\}$. By Remark 5.9, it is enough to show that $\text{tp}(\eta_0/B) = \text{tp}(\eta_N/B)$. By Corollary 5.15, we may assume without loss of generality that $N = n + 1$. It follows easily from Proposition 5.14 that for every $i \in (0, N)$, $i \neq u$ we have $\eta_0 \neq \eta_0 \wedge \eta_i = \eta_u \wedge \eta_i = \eta_N \wedge \eta_i \neq \eta_N$. Likewise, $\eta_0 \neq \eta_0 \wedge \eta_u = \eta_u \wedge \eta_{2u} = \eta_u \wedge \eta_N \neq \eta_N$. Using this, the conclusion follows easily from Fact 2.24. \square

The preceding corollary suggests the following question.

Question 5.17. If p is a finite partial automorphism of a tree, and η is its quasi-cyclic orbit, then does p admit an extension \bar{p} (to a finite partial automorphism, possibly of a larger tree) such that the orbit $\bar{\eta}$ of \bar{p} containing η is no longer quasi-cyclic?

Corollary 5.16 gives a positive answer under the additional assumption that p has no other orbits besides η . In general, this seems plausible, but we do not know the answer.

Remark 5.18. We believe that Corollary 5.16 can be extended to say the following (with a similar proof). If $\eta = (\eta_0, \dots, \eta_m)$ is a quasi-cycle with pseudo-period u , while $k > n$ is a multiple of u , then η can be extended to exactly the following kinds of finite partial orbits:

- a k -cycle,
- a quasi-cycle with pseudo-period u (of arbitrary length),
- a quasi-cycle with pseudo-period k (of arbitrary length),
- both an ascending and a descending spiral of spiral length k (of arbitrary length),
- both an ascending and a descending spiral comb of spiral length k (of arbitrary length),
- if $n < 2u$, both an ascending and a descending spiral comb of spiral length u (of arbitrary length). \diamond

6 Finding determined automorphisms

Recall that \mathbf{T} is the universal countable meet-tree.

Definition 6.1. Given a finite partial automorphism p , we call a point $a \in \text{dom}(p)$ an *initial point of p* if $a \notin \text{range}(p)$ or a is in a cyclic orbit of p . \diamond

Definition 6.2. Given sets P, Q , a partial function $p: P \rightarrow Q$, an element $a \in P$ and an integer k , we write $p^k(a)\downarrow$ if $a \in \text{dom}(p^k)$, i.e. if $p^k(a)$ is well-defined. Otherwise, we write $p^k(a)\uparrow$ (if $p^k(a)$ is undefined).

When p is fixed in the context, and we have an orbit $\eta = (\eta_0, \dots, \eta_m)$ of p , then we write also write $\eta_m\downarrow$ for $p^m(\eta_0)\downarrow$ and $\eta_m\uparrow$ for $p^m(\eta_0)\uparrow$. \diamond

Definition 6.3. Let p be a finite partial automorphism of \mathbf{T} . We say that it is *pseudo existentially closed (PEC)* (in \mathbf{T}) if it satisfies the following condition.

For every extension $f \supseteq p$ to a partial automorphism of \mathbf{T} , every triple η_0, μ_0, ζ_0 of initial points of p (cf. Definition 6.1) and every $m_1, m_2 > 0$ such that $f^{m_1}(\mu_0)\downarrow, f^{m_2}(\zeta_0)\downarrow$, there exist positive integers m'_1, m'_2, m''_1, m''_2 such that:

- if $\mu_{m_1}\uparrow$, then $m'_1 \neq m''_1$,
- the triples $(\eta_0, f^{m_1}(\mu_0), f^{m_2}(\zeta_0))$, $(\eta_0, \mu_{m'_1}, \zeta_{m'_2})$ and $(\eta_0, \mu_{m''_1}, \zeta_{m''_2})$ are well-defined, and all have the same quantifier-free type (in \mathbf{T} ; in particular, they have the same order type),
- if there is some k such that $\mu_0 \wedge f^k(\mu_0) \neq f^k(\mu_0) \wedge f^{2k}(\mu_0)$ (in particular, $f^{2k}(\mu_0)\downarrow$), then for minimal such k , we have $m_1 \equiv m'_1 \equiv m''_1 \pmod{k}$. \diamond

Proposition 6.4. *If p is pseudo existentially closed (PEC) and η, μ are its orbits, while $f \supseteq p$ is contained in some $\bar{f} \in \text{Aut}(\mathbf{T})$ (in particular, if f is a finite partial automorphism of \mathbf{T}), then:*

- (1) *if the f -orbit containing η is neither a quasi-cycle nor a cycle, then there is some k such that $\eta_{2k}\downarrow$ and $\eta_0 \wedge \eta_k \neq \eta_k \wedge \eta_{2k}$,*
- (2) *if the f -orbit containing η is a k -cycle, k -spiral (ascending or descending), k -comb (ascending or descending) or a quasi-cycle (respectively), then so is η (respectively),*
- (3) *if η, μ are distinct, then so are the f -orbits extending them.*

Proof. We may assume without loss of generality that $f = \bar{f} \in \text{Aut}(\mathbf{T})$. Write $\bar{\eta}$ for the f -orbit extending η , enumerated so that $\eta_0 = \bar{\eta}_0$ and $f(\bar{\eta}_i) = \bar{\eta}_{i+1}$, and likewise, let $\bar{\mu} \supseteq \mu$ be the f -orbit containing μ .

For (1), since $\bar{\eta}$ is not a quasi-cycle, there is some k such that $\eta_0 \wedge \bar{\eta}_k \neq \bar{\eta}_k \wedge \bar{\eta}_{2k}$. We may assume without loss of generality that this k is minimal. We claim that $\eta_{2k} \downarrow$, whence $\bar{\eta}_k = \eta_k$ and $\bar{\eta}_{2k} = \eta_{2k}$ and we are done. Otherwise, suppose towards contradiction that $\eta_{2k} \uparrow$, and apply PEC with $m_1 = 2k$, $m_2 = k$ and conclude that there positive m'_1, m'_2, m''_1, m''_2 with $m'_1 \equiv m''_1 \equiv 2k \pmod{k}$ such that $\eta_{m'_1} \downarrow, \eta_{m''_1} \downarrow$. But since m'_1, m''_1 are distinct and positive, it follows that $\max(m'_1, m''_1) \geq 2k$, so $\eta_{2k} \downarrow$, a contradiction.

For (2), suppose first that $\bar{\eta}$ is a k -spiral or a k -cycle, so η_0 is comparable to $\bar{\eta}_k$. We claim that $\eta_k \downarrow$, whence $\bar{\eta}_k = \eta_k$ and we are done. Indeed, by PEC (with $m_1 = m_2 = k$), there is some $m'_1 > 0$ such that $\eta_{m'_1} \downarrow$ and $\text{otp}(\eta_0, \eta_{m'_1}) = \text{otp}(\eta_0, \bar{\eta}_k)$. But then $\eta_{m'_1} = \bar{\eta}_{m'_1}$ is comparable to $\eta_0 = \bar{\eta}_0$, so by Proposition 5.6 k divides $m'_1 = m'_1 - 0$, whence $m'_1 \geq k$, so $\eta_k \downarrow$.

Otherwise, if $\bar{\eta}$ is a k -comb, then it is not a quasi-cycle, so η is not a quasi-cycle by (1). Since $\bar{\eta}$ is not a cycle, nor a spiral, neither is η , so it is an l -comb, whence $\bar{\eta}$ is an l -comb, so $l = k$.

Finally, if $\bar{\eta}$ is a quasi-cycle, then by Remark 5.4, so is η , completing the proof of (2).

Finally, (3) is immediate by PEC: if $\bar{\eta}_{m_1} = \bar{\mu}_{m_2}$, then we have some m'_1, m'_2 such that $\eta_{m'_1} = \mu_{m'_2}$. \square

Remark 6.5. Note that if the answer to Question 5.17 is positive, then it follows by Proposition 6.4(2) that a PEC partial automorphism has no quasi-cyclic orbits. This would make the proof of Theorem 6.6 below a bit simpler — namely, we could drop the final three paragraphs of the proof, and we would no longer need Proposition 5.14. \diamond

The following Theorem is one of the main building blocks of the proof that the universal countable meet-tree has a generic automorphism.

Theorem 6.6. *Let p be a PEC finite partial automorphism of \mathbf{T} .*

Let $\xi = (\xi_0, \xi_1, \dots, \xi_n)$ be an orbit of p such that no non-cyclic orbit of p has smaller length.

Then if $p \cup \{(\xi_n, v)\}$ and $p \cup \{(\xi_n, w)\}$ are both extensions of p to a partial automorphism of \mathbf{T} , and A is the meet-tree generated by the union of all orbits of p , then $\text{tp}(v/A) = \text{tp}(w/A)$.

Proof. Put $B := \langle \text{range}(p) \rangle$, $p_v := p \cup \{(\xi_n, v)\}$ and $p_w := p \cup \{(\xi_n, w)\}$. Note that the conclusion is trivially true if ξ is a cycle, so in the rest of the proof, we assume that ξ is not a cycle, and so $\xi_{n+1} \uparrow$. Furthermore, $\text{tp}(v/\text{range}(p)) = p(\text{tp}(\xi_n/\text{dom}(p))) = \text{tp}(w/\text{range}(p))$, and hence $\text{tp}(v/B) = \text{tp}(w/B)$.

Claim 1. *For every $\eta_i \in \text{dom}(p)$, we have $\text{tp}(\eta_i/B) \neq \text{tp}(v/B)$.*

Proof. Suppose towards contradiction that $\text{tp}(\eta_i/B) = \text{tp}(v/B)$. First, we show that in this case $p \cup \{(\xi_n, \eta_i)\}$ is a partial automorphism. Let \bar{c} enumerate $\text{dom}(p)$. Let $\varphi(x, \bar{c})$ be a quantifier-free formula. We need to show that if $\varphi(\xi_n, \bar{c})$ holds, then so does $\varphi(\eta_i, p(\bar{c}))$. But since p_v is a partial automorphism, $\varphi(\xi_n, \bar{c})$ holds if and only if $\varphi(v, p(\bar{c}))$ holds. But this is true if and only if $\varphi(x, p(\bar{c})) \in \text{tp}(v/B) = \text{tp}(\eta_i/B)$, i.e. $\varphi(\eta_i, p(\bar{c}))$ holds.

Now, if $\eta = \xi$ and $p \cup \{(\xi_n, \eta_i)\}$ is a partial automorphism, it means that ξ can be turned into a cycle in an extension of p . Since we have assumed ξ is not already a cycle, this contradicts Proposition 6.4(2). Otherwise, if $\eta \neq \xi$ and $p \cup \{(\xi_n, \eta_i)\}$ is a partial automorphism, we contradict Proposition 6.4(3) \square (claim)

Let $v' := \max_{b \in B} b \wedge v$ and likewise, $w' := \max_{b \in B} b \wedge w$.

Claim 2. $\text{tp}(v/A) = \text{tp}(w/A)$ if and only if the order type of v' over A is the same as the order type of w' over A (or some interval in $A \cup \{-\infty, \infty\}$, containing v' and w').

Proof. We have that $v' = \max_{a \in A} a \wedge v$. Indeed, otherwise, we would have $a \wedge v > v'$ for some $a \in A \setminus B$, and hence also for some $\eta_0 \in \text{dom}(p) \setminus \text{range}(p)$ (namely, any such that $\eta_0 \geq a$), we would have $\eta_0 \wedge v > v'$, and so, by Corollary 5.8, $\text{tp}(\eta_0/B) = \text{tp}(v/B)$, contradicting Claim 1. Analogously, $w' = \max_{a \in A} a \wedge w$.

Furthermore, $v = v'$ holds if and only if $w = w'$: they hold if and only if for some $b \in B$, the type $\text{tp}(v/B) = \text{tp}(w/B)$ implies $b \geq x$. Thus, the conclusion follows by Fact 2.24. \square (claim)

We divide the proof into two main cases, depending on whether $v' \in B$.

Case 1: $v' \in B$. In this case, we have in $\text{tp}(v/B)$ a formula $\varphi(x, B)$ which says “ $x \geq v'$ and $v' = \max_{b \in B} x \wedge b$ ”. Since $\text{tp}(v/B) = \text{tp}(w/B)$, we conclude w also satisfies this formula, so $v' = w'$, and in particular, they have the same order type over A , so by Claim 2, $\text{tp}(v/A) = \text{tp}(w/A)$, completing the proof in Case 1.

Case 2: $v' \notin B$. Note that as we have seen in the proof of Claim 2, we have that $v = v'$ if and only if $w = w'$, and analogously, $v' \in B$ if and only if $w' \in B$. We will treat separately the subcases when $v = v' \notin B$ and $v \neq v' \notin B$, but first, we make some observations that apply to both of them.

Put $\alpha := \max\{b \in B \mid b \leq v'\}$ ($-\infty$ if the set is empty), and choose $\beta \in B$ to be minimal above v' , so $\beta \geq v'$ and $(v', \beta) \cap B = \emptyset$. Then $v' \in (\alpha, \beta)$ and $(\alpha, \beta) \cap B = \emptyset$. Note that $\beta \wedge v = v'$. Indeed, $\beta \wedge v \leq v'$ by definition of v' , and since $\beta \geq v'$ and $v \geq v'$, also $\beta \wedge v \geq v'$. Since $\text{tp}(v/B) = \text{tp}(w/B)$, it follows that $w' \in (\alpha, \beta)$ and $w' = w \wedge \beta$. Notice that analogously, for every $b \in B$ such that $b \geq \beta$, we have also $b \wedge v = v'$ and $b \wedge w = w'$.

In the analysis of the two subcases, we will use the α and β heavily.

Claim 3. Every element of $A \cap (\alpha, \beta)$ is of the form $\eta_0 \wedge \beta$ for some $\eta_0 \in \text{dom}(p) \setminus B$.

Proof. Let $a \in A \cap (\alpha, \beta)$. Then for some η, μ, i, j we have $a = \eta_i \wedge \mu_j$. Since $a < \beta$, we have $\eta_i \wedge \beta \geq a$ and $\mu_j \wedge \beta \geq a$. By Fact 2.18, only one of these inequalities can be strict. Suppose without loss of generality that $\eta_i \wedge \beta = a$. Since $a \in (\alpha, \beta)$, it follows that $a \notin B$, so (because $\beta \in B$) $\eta_i \notin B$, so also $\eta_i \notin \text{range}(p)$, whence $\eta_i \in \text{dom}(p)$, which completes the proof. \square (claim)

Since $v', w' \in (\alpha, \beta)$, by Claim 3 and Claim 2, it is enough to show that for every $\eta_0 \in \text{dom}(p) \setminus B$, at least two of the following hold:

- $\eta_0 \wedge \beta \leq v'$ if and only if $\eta_0 \wedge \beta \leq w'$,
- $\eta_0 \wedge \beta \geq v'$ if and only if $\eta_0 \wedge \beta \geq w'$,
- $\eta_0 \wedge \beta = v'$ if and only if $\eta_0 \wedge \beta = w'$.

(Because any two of these conditions imply that v', w' have the same order type over the interval (α, β) .)

Now, if $\eta_0 \wedge \beta \notin (\alpha, \beta)$, then it is clear that these equivalences hold (because $v', w' \in (\alpha, \beta)$). Hence, for the rest of the proof, let us fix an arbitrary $\eta_0 \in \text{dom}(p) \setminus B$ such that $\eta_0 \wedge \beta \in (\alpha, \beta)$. The goal is to show that for this η_0 , two of the above three equivalences hold.

Case 2.1: $v' = v \notin B$.

If $\eta_0 \wedge \beta \geq v'$, then trivially $\eta_0 \geq v$. Since p_v extends p (as a partial automorphism of \mathbf{T}), $v = p_v^{n+1}(\xi_0) \downarrow$ and $\xi_{n+1} \uparrow$, by Definition 6.3 (applied to $m_1 = m_2 = n + 1$, $f = p_v$), there are $i < j \leq n$ such that $\xi_i, \xi_j < \eta_0$, and hence they are comparable. Since ξ is not a cycle, it follows that it must be a spiral (see Definition 5.3 to recall the orbit types). Thus, if ξ is not a spiral, then for all $\eta_0 \in \text{dom}(p) \setminus B$ we have $\eta_0 \wedge \beta < v'$ and analogously, $\eta_0 \wedge \beta < w'$.

By the preceding paragraph, we may assume that ξ is spiral, of spiral length k for some k . We will show that $\eta_0 \wedge \beta \geq v$ (equivalently, $\eta_0 \geq v$) if and only if $\eta_0 \geq \xi_{n+1-k}$ (and so analogously $\eta_0 \wedge \beta \geq w$ if and only if $\eta_0 \geq \xi_{n+1-k}$).

Suppose first that ξ is descending. Then $v < \xi_{n+1-k}$, so if $\eta_0 \geq \xi_{n+1-k}$, then trivially $\eta_0 \geq v$. Conversely, if $\eta_0 \geq v$, then by Definition 6.3 (and the equality $v = p_v^{n+1}(\xi_0)$), for some $i \equiv n + 1 \pmod{k}$, we have $\eta_0 \geq \xi_i$. Since $i \equiv n + 1 \equiv n + 1 - k \pmod{k}$ and ξ is descending, $\xi_i \geq \xi_{n+1-k}$ and so $\xi_{n+1-k} \leq \eta_0$.

Otherwise, suppose ξ is ascending. Then $v > \xi_{n+1-k}$, so if $\eta_0 \geq v$, then trivially $\eta_0 \geq \xi_{n+1-k}$. Conversely, if $\eta_0 \not\geq v$, then again by Definition 6.3, there is some $i \equiv n + 1 \pmod{k}$, $i \leq n$, such that $\xi_i \not\leq \eta_0$, and as before, we have $\xi_i \leq \xi_{n+1-k}$, so $\xi_{n+1-k} \not\leq \eta_0$.

To complete the proof in this case, it is enough to show that given $\eta_0 \in \text{dom}(p) \setminus B$, we have $\eta_0 \wedge \beta = v$ if and only if $\eta_0 \wedge \beta = w$. We will show that these equalities are impossible (given that ξ is spiral and $v = v' \notin B$, and hence $w = w' \notin B$). Notice that for some $\mu_i \in \text{range}(p)$ (so without loss of generality, $i > 0$) we have that $\mu_i \geq \beta$. Since $\beta = \mu_i \wedge \beta > \beta \wedge \eta_0$, by Fact 2.18, it follows that $\beta \wedge \eta_0 = \mu_i \wedge \eta_0$.

Suppose ξ is a descending k -spiral and $\eta_0 \wedge \beta = v$. By Definition 6.3 (with $f = p_v$, $m_1 = n + 1$, $m_2 = i$), there is some $j \equiv n + 1 \pmod{k}$ and $i' > 0$ such that $\eta_0 \wedge \mu_{i'} = \xi_j$. Since ξ is descending, $\xi_j > v$, so $\eta_0 \wedge \mu_{i'} > \eta_0 \wedge \mu_i = v$, and by Fact 2.18, we have that $v = \eta_0 \wedge \mu_i = \mu_i \wedge \mu_{i'}$. Since $i, i' > 0$, it follows that $v \in B$, a contradiction.

Now, suppose ξ is an ascending k -spiral and $\eta_0 \wedge \beta = v$. Arguing as in the preceding paragraph, we conclude that for some $j \equiv n + 1 \pmod{k}$ and a positive i' such that $\eta_0 \wedge \mu_{i'} = \xi_j$. We may assume without loss of generality that i' is maximal (i.e. for all $i'' > i'$ there is no $j' \equiv n + 1 \pmod{k}$ such that $\eta_0 \wedge \mu_{i''} = \xi_{j'}$).

We claim that $i' \leq i - k$. Indeed, put $l := i' - i$ and let $f \in \text{Aut}(\mathbf{T})$ be an arbitrary extension of p_v ; note that the f -orbit of ξ_j is still an ascending k -spiral, and also that $f^k(\xi_j) \leq f^{n+1-j}(\xi_j) = v \leq \eta_0$, so in particular, $f^k(\xi_j) \not\leq \mu_{i'}$ — otherwise, $f^k(\xi_j) \leq \mu_{i'} \wedge \eta_0 = \xi_j$, which would contradict the assumption that the f -orbit of ξ_j is an ascending k -spiral. Since $\eta_0 \wedge \mu_i = v > \xi_j = \eta_0 \wedge \mu_{i'}$, we have $\mu_i \wedge \mu_{i'} = \xi_j$, so $\mu_{i'} \wedge f^l(\mu_{i'}) = f^l(\xi_j)$. Now, $\mu_{i'} \geq \xi_j$, $f^l(\xi_j)$, so the latter two are comparable, whence k divides l (by Proposition 5.6); this implies that $f^l(\xi_j)$ and $f^k(\xi_j)$ are comparable. On the other hand, since $f^k(\xi_j) \not\leq \mu_{i'}$ and $f^l(\xi_j) \leq \mu_{i'}$, whence $f^l(\xi_j) < f^k(\xi_j)$, so $l < k$. Since k divides l and l cannot be 0, it follows that $l \leq -k$.

It follows that in fact $j = n + 1 - k$. Otherwise, if $j \neq n + 1 - k$, then $\xi_{n+1-k} > \xi_j = \mu_i \wedge \mu_{i'}$. On the other hand, $\mu_{i-k} \geq \xi_{n+1-k}$, so $\mu_{i-k} \wedge \mu_i \geq \xi_{n+1-k}$, and hence $\mu_{i-k} \wedge \mu_{i'} = \xi_j$. It would follow that $\mu_k \wedge \mu_{i'+k} = \xi_{j+k}$. Now, $\mu_k \wedge \eta_0 = v > \xi_{j+k}$, so $\mu_{i'+k} \wedge \eta_0 = \xi_{j+k}$, contradicting the maximality of i' .

Thus, we have that $i' \leq i - k$ and $\mu_{i'} \wedge \eta_0 = \xi_{n+1-k}$. It follows that $\mu_{i'+k} \wedge \eta_k = v$ ($\mu_{i'+k}$ is well-defined because $i' \leq i - k$ and μ_i is well-defined, while η_k is well-defined because ξ is the shortest non-cyclic orbit). But $\mu_{i'+k} \wedge \eta_k \in B$, a contradiction.

Thus, we have completed the proof in Case 2.2 (i.e. under the assumption that $v = v' \notin B$). The following is the last remaining case.

Case 2.2: $v' < v$ and $v' \notin B$. Recall that we have fixed some $\eta_0 \in \text{dom}(p) \setminus B$ such that $\eta_0 \wedge \beta \in (\alpha, \beta)$, and we need to show that it compares to v' in the same way as it compares

to w' .

Claim 4. *In the case we are considering, $\eta_0 \in (\alpha, \beta)$.*

Proof. We have that $\eta_0 \wedge \beta \in (\alpha, \beta)$. Let $\bar{b} \in B$ be such that $\eta_0 \wedge \bar{b} = \max_{b \in B} \eta_0 \wedge b$. Then $\eta_0 \wedge \beta \leq \eta_0 \wedge \bar{b}$. If the inequality is strict, then by Fact 2.18, we have $\eta_0 \wedge \beta = \bar{b} \wedge \beta \in B \cap (\alpha, \beta)$, which contradicts the definition of α and β . Otherwise, $\eta_0 \wedge \beta = \eta_0 \wedge \bar{b} \in (\alpha, \beta)$, so $\max_{b \in B} \eta_0 \wedge b$ has the same order type over the interval (α, β) as v . Since, by Claim 1, $\text{tp}(v/B) \neq \text{tp}(\eta_0/B)$, by Fact 2.24 and the inequality $v \neq v'$, we conclude that $\eta_0 = \max_{b \in B} \eta_0 \wedge b = \eta_0 \wedge \beta$, and therefore $\eta_0 = \eta_0 \wedge \beta \in (\alpha, \beta)$. \square (claim)

Claim 5. *We have that $\eta_0 = v \wedge \mu_k$ for some orbit μ and $k > 0$ if and only if there are $l, r > 0$ such that $\xi_n \wedge \mu_{k-1} = \xi_{l-1} \wedge \mu_{r-1}$ and $\eta_0 = \xi_l \wedge \mu_r$.*

In particular, $\eta_0 = v'$ if and only if $\eta_0 = w'$.

Proof. Suppose $\eta_0 = v \wedge \mu_k$. By Definition 6.3 (applied to $f = p_v$, $m_1 = n + 1$, $m_2 = k$), we have some l, r such that $\eta_0 = \xi_l \wedge \mu_r$, and they clearly satisfy the right-hand side. The converse is immediate (just apply p_v to both sides of the first equality).

For “in particular”, just note that $v' = v \wedge \mu_k$ for some μ, k (namely, any such that $\mu_k \in B$ and $\mu_k \geq \beta$) and $w' = w \wedge \mu_k$ for the same μ, k (because $\text{tp}(v/B) = \text{tp}(w/B)$) and apply the first part. \square (claim)

By Claim 4 and Claim 5, it is enough to show that we have $\eta_0 \leq v'$ if and only if $\eta_0 \leq w'$. Note that since $v' \in (\alpha, \beta)$, this implies that v' and η_0 are comparable. Note also that since $v' < v$ and $\eta_0 < \beta$, it is easy to see that $\eta_0 \leq v'$ if and only if $\eta_0 < v$.

Suppose ξ is not a quasi-cycle, by Proposition 6.4(1), for some (minimal) k , we have that $\xi_0 \wedge \xi_k \neq \xi_k \wedge \xi_{2k}$ (so by semilinearity, there is a strict inequality). We claim that under this assumption, $\eta_0 \leq v'$ (since $\eta_0 < \beta$, equivalently, $\eta_0 \leq v$) if and only if $\eta_0 \leq \xi_{n+1-k} \wedge \xi_{n+1-2k}$.

If $\xi_0 \wedge \xi_k < \xi_k \wedge \xi_{2k}$, then by Proposition 5.7, $\xi_{n+1-k} \wedge \xi_{n+1-2k} < \xi_{n+1-k} \wedge v$, which makes one implication trivial. In the other direction, if $v \geq \eta_0$, then since $\xi_{n+1} \uparrow$, by Definition 6.3 (applied to $f = p_v$, $m_1 = m_2 = n + 1$), there are distinct $i_1, i_2 \equiv n + 1 \pmod{k}$ such that $\xi_{i_1}, \xi_{i_2} \geq \eta_0$, whence $\xi_{i_1} \wedge \xi_{i_2} \geq \eta_0$. We may assume without loss of generality that $i_1 < i_2$, whence $i_1 \leq n + 1 - 2k$ (because $i_1 \equiv i_2 \equiv n + 1 \pmod{k}$ and $i_2 \leq n$). By Proposition 5.7, it follows that $\xi_{i_1} \wedge \xi_{i_2} \leq \xi_{n+1-2k} \wedge \xi_{n+1-k}$, and hence $\eta_0 \leq \xi_{n+1-2k} \wedge \xi_{n+1-k}$.

Otherwise, suppose $\xi_0 \wedge \xi_k > \xi_k \wedge \xi_{2k}$. If $\eta_0 \leq v$, then by Definition 6.3 (with $f = p_v$, $m_1 = m_2 = n + 1$), there is some $i \equiv n + 1 \pmod{k}$ such that $\xi_i \geq \eta_0$, whence $\xi_i \wedge v \geq \eta_0$. Then by Proposition 5.7, $\xi_i \wedge v = \xi_{n+1-k} \wedge v < \xi_{n+1-2k} \wedge \xi_{n+1-k}$. Conversely, suppose $\eta_0 \not\leq v$. Then, again by Definition 6.3 (having in mind that $\xi_{n+1} \uparrow$), there are distinct $i_1, i_2 \equiv n + 1 \pmod{k}$ such that $\eta_0 \not\leq \xi_{i_1}, \xi_{i_2}$, so in particular, $\eta_0 \not\leq \xi_{i_1} \wedge \xi_{i_2}$; since $i_1, i_2 \leq n$ and $i_1, i_2 \equiv n + 1 \pmod{k}$, we have $\max(i_1, i_2) \leq n + 1 - k$, so $\min(i_1, i_2) \leq n + 1 - 2k$, and thus by Proposition 5.7, $\xi_{i_1} \wedge \xi_{i_2} \geq \xi_{n+1-k} \wedge \xi_{n+1-2k}$, whence $\xi_{n+1-k} \wedge \xi_{n+1-2k} \not\leq \eta_0$.

We are left with the case when ξ is a quasi-cycle. Note that by Proposition 6.4(2), it follows that ξv is also a quasi-cycle. Let u be the pseudo-period of ξ (cf. Definition 5.10). We will show that $v > \eta_0$ if and only if there is some positive $k \neq n + 1 - u$ such that $\xi_k \wedge \xi_{n+1-u} \geq \eta_0$.

Suppose $v > \eta_0$. Considering $\xi_{n+1} \uparrow$, by Definition 6.3, there are distinct positive $k, k' \leq n$ such that $\xi_k, \xi_{k'} > \eta_0$. Note that this implies that $v \wedge \xi_k \geq \eta_0$, so also $v \wedge \xi_{n+1-u} \geq \eta_0$, and in particular, $\xi_{n+1-u} \geq \eta_0$. Since $k \neq k'$, we may assume without loss of generality that $k \neq n + 1 - u$, and then clearly $\xi_k \wedge \xi_{n+1-u} \geq \eta_0$.

Now, suppose $0 < k \neq n + 1 - u$ and $\xi_k \wedge \xi_{n+1-u} \geq \eta_0$. Then by Proposition 5.14, $\xi_n \wedge \xi_{k-1} = \xi_{n-u} \wedge \xi_{k-1}$, so also $v \wedge \xi_k = \xi_{n+1-u} \wedge \xi_k \geq \eta_0$, and hence $v > \eta_0$.

This finishes the proof in Case 2.2, completing the proof of Theorem 6.6. \square

Remark 6.7. It is easy to see that an immediate extension (in the sense of Definition 4.7) of a PEC automorphism of \mathbf{T} is itself PEC. \diamond

Corollary 6.8. *A PEC partial automorphism of \mathbf{T} is determined (in the sense of Definition 4.1).*

Proof. It is enough to check that a PEC partial automorphism satisfies the hypothesis of Lemma 4.10. Take $p = f_0 \subseteq f_1 \subseteq \dots \subseteq f_n \subseteq f_{n+1}$ is as in the hypothesis and a positively strict extension $g \supseteq f_n$. Let $a = \nu_m$ be the sole element of $\text{dom}(f_{n+1}) \setminus \text{dom}(f_n)$. Put $v := f_{n+1}(a), w := g(a)$. Then by the preceding remark, f_n is PEC, so by Theorem 6.6, $\text{tp}(v/\text{dom}(f_n)) = \text{tp}(w/\text{dom}(f_n))$, so, by ultrahomogeneity of \mathbf{T} , there is some $\sigma \in \text{Aut}(\mathbf{T}/\text{dom}(f_n))$ such that $\sigma(v) = w$. It follows that $\sigma \circ f_n \circ \sigma^{-1}(f_{n+1}) \subseteq g$, and so $\tau_n := \sigma \upharpoonright_{\text{dom}(f_{n+1})}$ witnesses that the hypothesis of Lemma 4.10 is satisfied, which completes the proof. \square

Proposition 6.9. *Every finite partial automorphism of \mathbf{T} can be extended to a finite PEC partial automorphism.*

Proof. Take any partial automorphism p . We extend it in three steps:

1. Given an orbit η of p , if in some extension $f \supseteq p$, η is extended to an orbit which is not a quasi-cycle, then we extend η so that it is witnessed already in p ; we repeat that for each orbit of p .
2. For each triple η_0, μ_0, ζ_0 of initial points of p , each residue mod the spiral length k of μ (if it exists) and each quantifier-free type $r(x, y, z)$, we check whether for some $f \supseteq p$ there are m_1, m_2 (with m_1 having the appropriate residue mod k) such that $\models r(\eta_0, f^{m_1}(\mu_0), f^{m_2}(\zeta_0))$. If the answer is yes, we extend η, μ and ζ to already contain witnesses $\mu_{m'_1}, \zeta_{m'_2}$ for that.
3. We repeat the previous step, only this time, checking whether it is possible to have two witnesses $(\mu_{m'_1}, \zeta_{m'_2})$ and $(\mu_{m''_1}, \zeta_{m''_2})$ with $m'_1 \neq m''_1$.

This procedure terminates: we never add any new orbits in any of the steps, so step 1 completes. After step 1, the set of initial points is fixed (because there are no new orbits, and step 1 ensures that there can be no new cycles), as are the spiral lengths of all the orbits, so in each of steps 2 and 3 there is only a finite number of conditions to check.

It is fairly easy to check that after the three steps, we obtain a PEC partial automorphism. Step 1 guarantees that the minimal k checked in the last bullet of Definition 6.3 is simply the spiral length of μ . Step 2 provides the witness (m'_1, m'_2) required in the second bullet. Finally, at this point, the hypothesis of the first bullet of Definition 6.3 implies that $m_1 \neq m'_1$, so by the third step, we also have (m''_1, m''_2) with $m'_1 \neq m''_1$ (otherwise, we can just take $m''_1 = m'_1$ and $m''_2 = m'_2$). \square

Corollary 6.10. *The class \mathcal{K}_p^1 of finite meet-trees with a single partial automorphism has CAP and JEP.*

Proof. By Corollary 6.8 and Corollary 4.18, if p is a PEC automorphism of \mathbf{T} , then (B, p) is an amalgamation base in \mathcal{K}_p^1 , where $B \subseteq \mathbf{T}$ is generated by $\text{dom}(p) \cup \text{range}(p)$. Given any $(B', p') \in \mathcal{K}_p^1$, we may assume that $B' \subseteq \mathbf{T}$, and then extend p' to a finite partial automorphism p'' of \mathbf{T} such that $B' \subseteq \text{dom}(p'')$. and then by Proposition 6.9, we can extend p'' to a PEC partial automorphism p of \mathbf{T} , yielding a B as above such that $(B', p') \subseteq (B, p)$, which shows CAP.

To see JEP, take any $(A, p_A), (B, p_B) \in \mathcal{K}_p^1$. We may assume without loss of generality that $A \cap B = \emptyset$. Let v be a new element. Then $A \vee_v B := (A \cup B \cup \{v\}, p_A \cup p_B) \in \mathcal{K}_p^1$, ordered in such a way $a \wedge b = v$ for all $a \in A$ and $b \in B$, and clearly (A, p_A) and (B, p_B) embed into $A \vee_v B$. \square

The following theorem, along with Corollary 3.8, completes the **Main Theorem**.

Theorem 6.11. *The universal meet-tree \mathbf{T} has a generic automorphism.*

Proof. This is immediate by Corollary 6.10 and Fact 2.9. \square

Remark 6.12. It seems like it might be possible to use largely the same methods to show that for each $k > 0$, \mathbf{T}_k has a generic automorphism. More specifically, we believe that the analogue of Theorem 6.6 is true with essentially the same proof (after we replace the notion of PEC by the appropriate variant for \mathbf{T}_k), which will yield a cofinal class of *determined* automorphisms of \mathbf{T}_k . On the other hand, Proposition 4.16 shows that the analogue of Lemma 4.15 fails if $k > 1$, so we cannot simply use Lemma 4.3 to conclude that the class of determined automorphisms witnesses CAP, as we did in the proof of Corollary 6.10.

Instead, one could try to show that every finite partial automorphism of \mathbf{T}_k can be extended to one which is PEC (in \mathbf{T}_k) and has the property that its unique strict extension satisfies the hypothesis of Remark 4.17, which (together with Lemma 4.3 and a proof of Remark 4.17) would yield CAP. For $k = 2$, this seems straightforward, but the general case appears to be more difficult — it is plausible that for $k > 2$ we might have EAP but not CAP, leading to a negative answer to Question 1.2.

Note that in any case we easily get JEP for the class of finite k -ary meet-trees for all k : for $k > 1$ it is exactly as in the proof of Corollary 6.10, and for $k = 1$ (i.e., linear orders) it is also quite easy. Thus, for all k , the automorphism group of \mathbf{T}_k has a dense conjugacy class. \diamond

Remark 6.13. The caveats mentioned in Remark 6.12 do not seem to apply in the case of $k = 1$ (see Remark 4.20), so we can slightly adjust the proof of Theorem 6.11 to recover the fact that $(\mathbf{Q}, <)$ has a generic automorphism. \diamond

Remark 6.14. Consider the class of finite meet-trees expanded with a lexicographic ordering (i.e. a total order \triangleleft , extending the tree order, such that if $b \triangleleft a$ and $a \wedge a' > a \wedge b$, then $b \triangleleft a'$). It is easy to see that it is a Fraïssé class, and one can ask whether the limit of this class has a generic automorphism.

The proof of Lemma 4.15 seems to adapt to this context in a straightforward manner, so the main difficulty seems to lie with the analogue of Theorem 6.6.

Note that the orbit analysis for trees with a lexicographic ordering is much simpler: there are no nontrivial cycles (there can be fixed points), which implies that the spiral length can only be equal to 1 (so the only spirals are just monotone sequences, the only spiral combs are 1-combs), and the pseudo-period of a quasi-cycle is always one (so all quasi-cycles are “fans”). \diamond

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ITAY KAPLAN

The Hebrew University of Jerusalem

 <https://orcid.org/0000-0002-7032-1710>

kaplan@math.huji.ac.il

TOMASZ RZEPECKI

The Hebrew University of Jerusalem and Uniwersytet Wrocławski

 <https://orcid.org/0000-0001-9786-1648>

tomasz.rzepecki@math.uni.wroc.pl

DAOUD SINIORA



daoud.siniora@gmail.com